

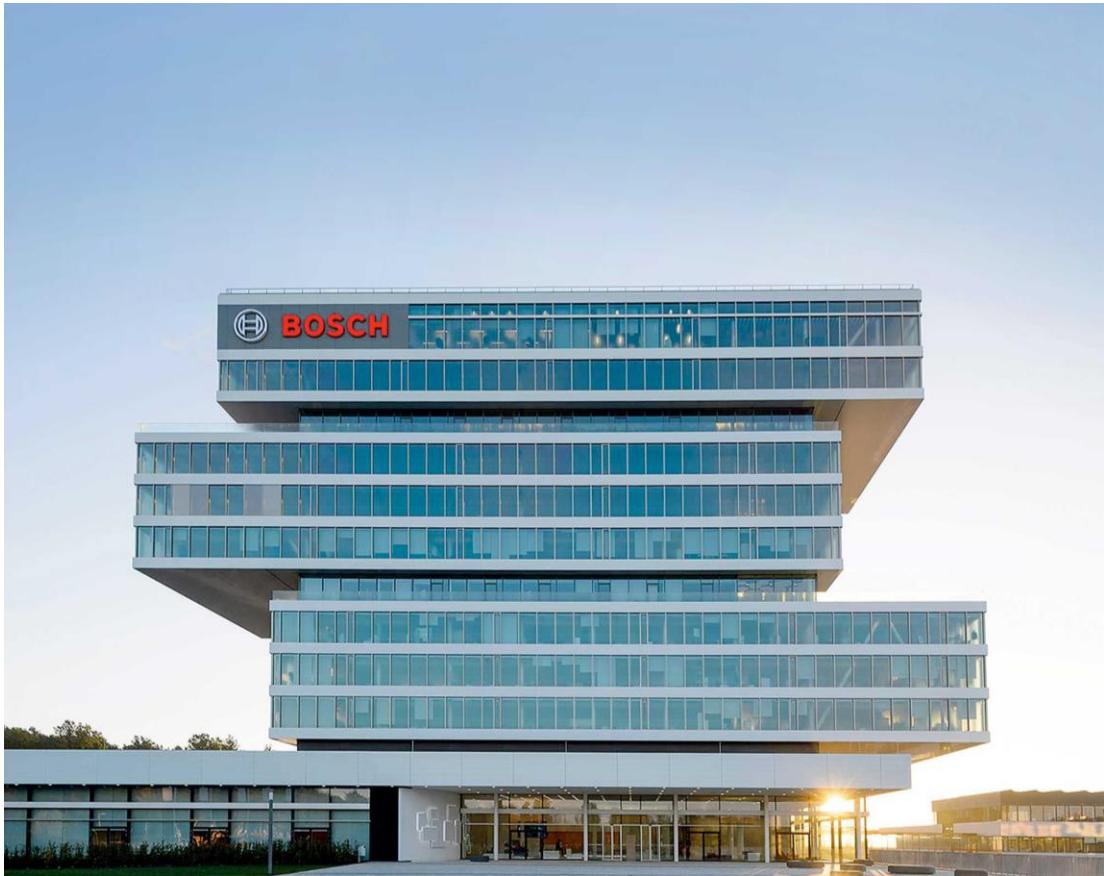
# SYNCHRONOUS LANGUAGES

## LECTURE 14 THE BLECH PROGRAMMING LANGUAGE

16 JUNE, 2020  
FRIEDRICH GRETZ  
BOSCH CORPORATE RESEARCH



# Today's speaker



Dr. Friedrich Gretz  
Robert Bosch GmbH  
Corporate Research in Renningen

[Friedrich.Gretz@de.bosch.com](mailto:Friedrich.Gretz@de.bosch.com)  
[www.blech-lang.org](http://www.blech-lang.org)

# Overview

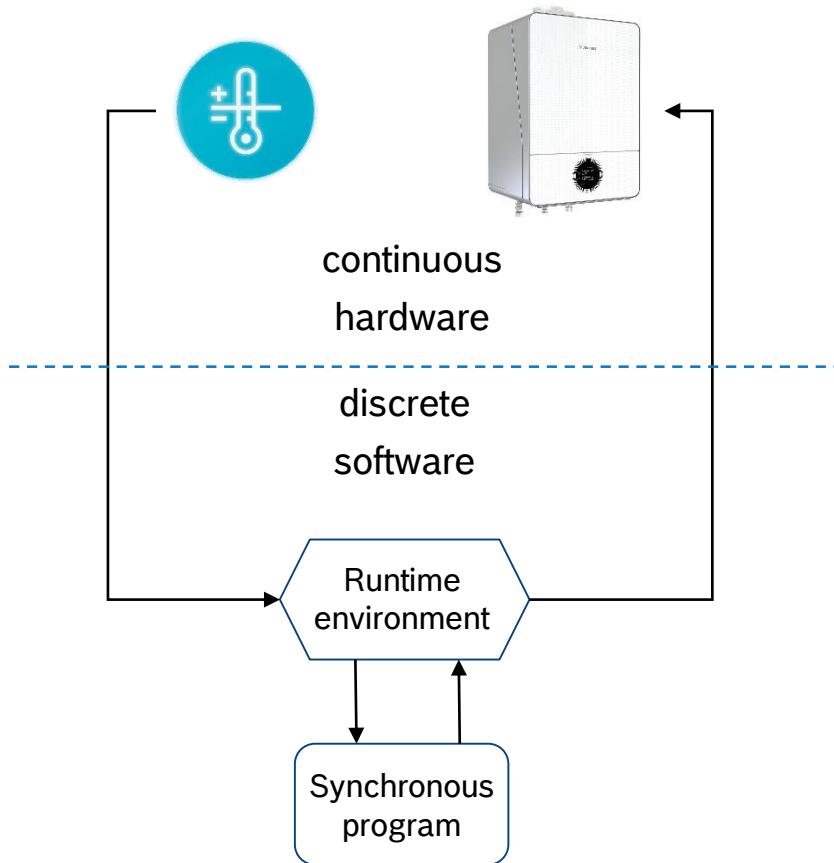
- ▶ Today's speaker
- ▶ **Why is synchronous programming interesting for Bosch?**
- ▶ Design goals
- ▶ Blech – as of now
- ▶ Application examples
- ▶ Outlook on planned features
- ▶ Additional remarks

# Why is synchronous programming interesting for Bosch? Reactive, embedded software everywhere!



# Abstract view of a reactive system

## Where do we use a synchronous language?



- ▶ Environment communicates asynchronously with physical world, drives synchronous programs
- ▶ A program is executed in *steps*
  - Assume a step takes no time (happens instantaneously)
  - No change of input data throughout computation
- ▶ A sequence of steps is called a *thread of execution*
- ▶ Threads can be composed concurrently
  - Accesses to shared data happen in a deterministic, causal order

# Do we need a new synchronous language? Available alternatives do not fulfill our requirements

- ▶ Céu              purely event-triggered, no causality, soft-realtime
- ▶ Esterel            no longer supported, not sequentially constructive, not separately compilable
- ▶ Lustre            not imperative, good for evaluating control loop equations but less intuitive for describing step-wise, mode switching behaviour
- ▶ SCCharts        automata centric view

Create a synchronous imperative language – Blech

# Overview

- ▶ Today's speaker
- ▶ Why is synchronous programming interesting for Bosch?
- ▶ **Design goals**
- ▶ Blech – as of now
- ▶ Application examples
- ▶ Outlook on planned features
- ▶ Additional remarks

# Design goals

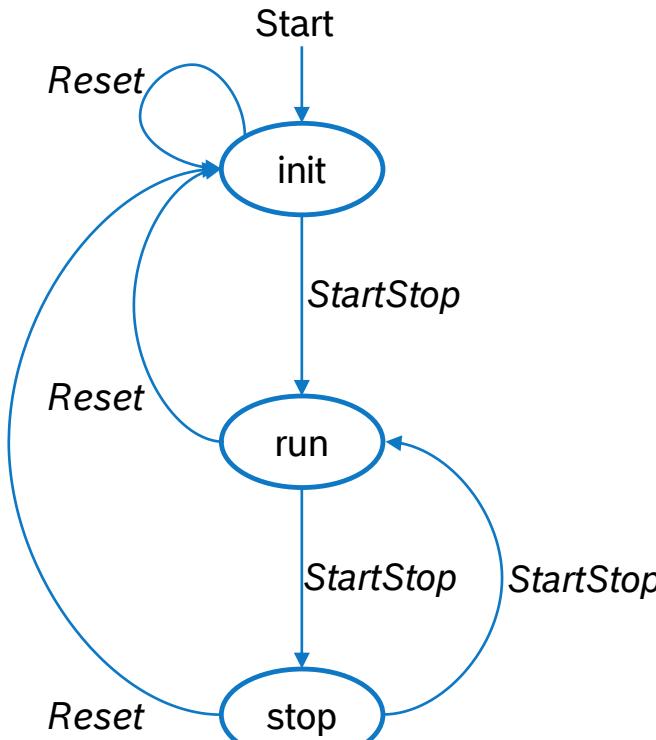
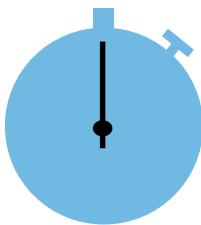
## Requirements

- ▶ Clear focus
  - ▶ Software
  - ▶ Reactive
  - ▶ Resource-constrained
  - ▶ Real-time
  - ▶ Scalable
- ▶ Deployment
  - ▶ Efficient code generation
  - ▶ Safe code generation
  - ▶ Integrate synchronous “execution shell” with existing real-time OS environments
  - ▶ Deployment on multi-core platforms
- ▶ Domain orientation
  - ▶ Embedded
  - ▶ Control intensive systems
  - ▶ Computations and switching behaviour
  - ▶ Intertwined functionality
- ▶ Developer Orientation
  - ▶ Readable
  - ▶ Clear semantics
  - ▶ Stateflow in controlflow
  - ▶ Structured data
  - ▶ Code structuring, information hiding
  - ▶ Safe and modern type system
- ▶ Compatibility
  - ▶ Integration **of** legacy code
  - ▶ Integration **in** legacy code
  - ▶ Support separate compilation
- ▶ Testing & Safety
  - ▶ Deterministic, repeatable testing
  - ▶ Integrate with existing simulation frameworks
  - ▶ Reduce false positives in static code analysis
  - ▶ Provide more guarantees, e.g. through causality

# Overview

- ▶ Today's speaker
- ▶ Why is synchronous programming interesting for Bosch?
- ▶ Design goals
- ▶ **Blech – as of now**
- ▶ Application examples
- ▶ Outlook on planned features
- ▶ Additional remarks

# Mode transitions as synchronous control flow



```

activity stopWatchControl (isPressedStartStop: bool,
                           isPressedReset: bool)
                           (display: Display)

when isPressedReset reset
  // init
  resetToZero()(display)
  if not isPressedStartStop then
    await isPressedStartStop
  end
  repeat
    // run
    repeat
      await true
      increment()(display)
    until isPressedStartStop end
    // stop
    await isPressedStartStop
  end
end
end
  
```

## Concurrent composition of behaviours over time

```
// Main Program
@[EntryPoint]
activity Main (isPressedStartStop: bool,
                isPressedReset: bool)
    var display: Display
    cobegin // render
        repeat
            show(display)
            await true
        end
    with // control
        run StopWatchController(isPressedStartStop,
                                isPressedReset)
                                (display)
    end
end
```

- ▶ Execution model
  - ▶ Concurrent behaviours run in synchronised steps
- ▶ Causal order
  - ▶ first, update display data
  - ▶ second, show display
- ▶ Code generation
  - ▶ sequential code
  - ▶ Statically ordered by the compiler

# Blech

## Concurrency in detail

```
cobegin [weak]
    ...
        do a step here
    and
    ...
        do a step there
end
```

```
cobegin
    run A(x)(z)
with
    run B(y)(z)
end
```

write-write conflict  
reject compilation

```
cobegin
    run A(x)(z)
with
    run B(y)(x)
end
```

in every reaction:  
“write before read!”

```
cobegin
    run A(x)(z)
with
    run B(z)(x)
end
```

write-read cycle  
reject compilation

```
cobegin
    run A(x)(z)
with
    run B(prev z)(x)
end
```

solution  
use previous value

# Blech

## Concurrency in detail

```
...  
cobegin  
...  
with  
...  
...  
cobegin  
    run A(x)(z)  
with  
...  
end  
...  
with  
    run B(prev z)(x)  
end  
...  
...
```

- Cobegin may have any fixed number of blocks
- Cobegin is orthogonal: it can be arbitrarily nested
- Subprograms are **black boxes** with interfaces, may be **compiled separately**
- **Interfaces** tell what data types are expected **and** whether data is only **read** or also **written**
- Causal **scheduling is dealt with locally** at call site
- Causality issues arise and may be debugged and **fixed within one cobegin statement!**

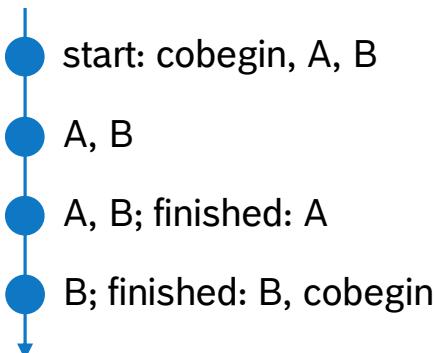
# Blech

## Concurrency in detail

```
cobegin  
    run A(x)(z)  
with  
    run B(y)(x)  
end
```

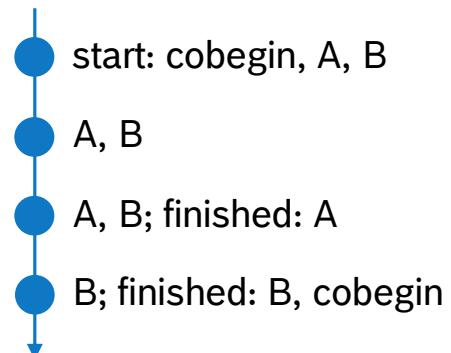
cobegin statement terminates when...

A **and** B have finished all their reactions



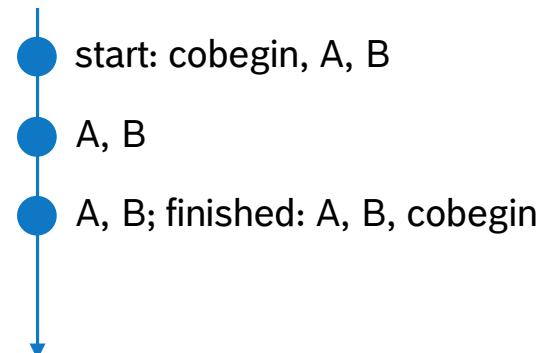
```
cobegin weak  
    run A(x)(z)  
with  
    run B(y)(x)  
end
```

B has finished all its reactions;  
A is possibly aborted



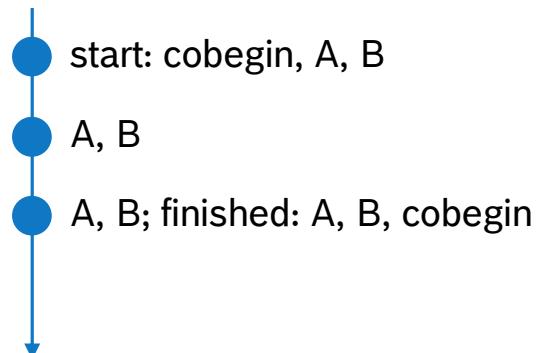
```
cobegin  
    run A(x)(z)  
with weak  
    run B(y)(x)  
end
```

A has finished all its reactions;  
B is possibly aborted



```
cobegin weak  
    run A(x)(z)  
with weak  
    run B(y)(x)  
end
```

A **or** B has finished all its reactions;  
the other one is possibly aborted



# Blech

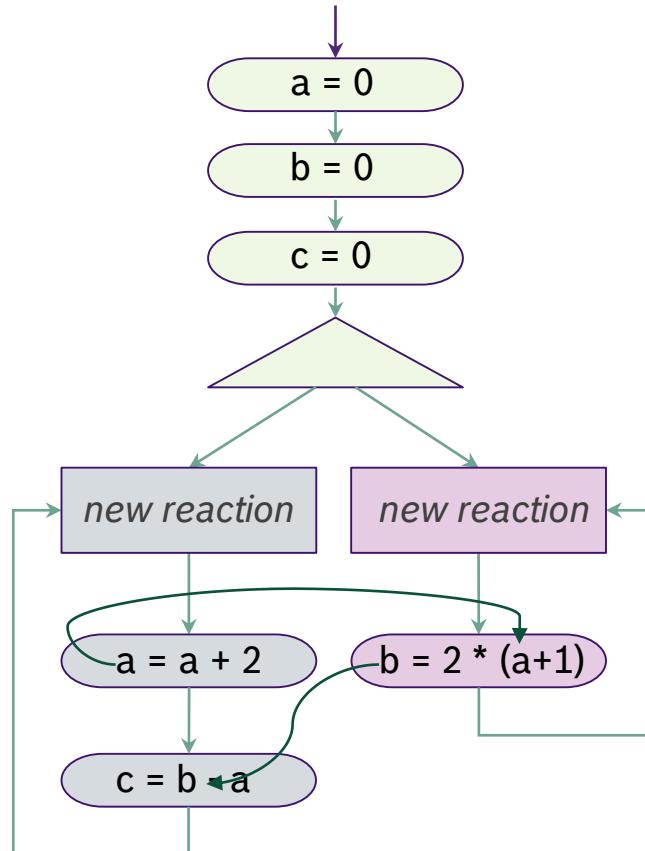
## Use case for weak branches

```
runs indefinitely, unless terminated  
cobegin weak  
    run BlinkLED(...)(...)          // no arguments for readability  
with  
    run WaitForKeyStroke(...)(...) // no arguments for readability  
end  
eventually terminates (if the system is to make any progress at all)
```

# Blech

## Compiling activities to sequential C functions

```
@[EntryPoint]
activity main()()
var a: int32
var b: int32
var c: int32
cobegin
    repeat
        await true
        a = a + 2
        c = b - a
    end
    with
        repeat
            await true
            b = 2 * (a + 1)
        end
    end
end
```



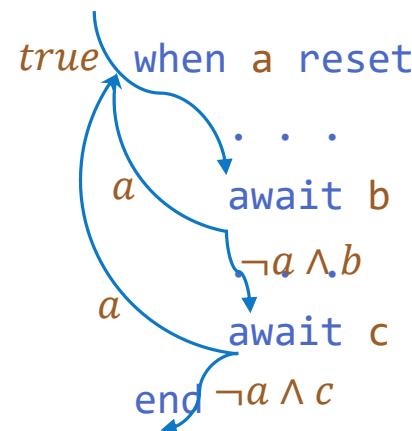
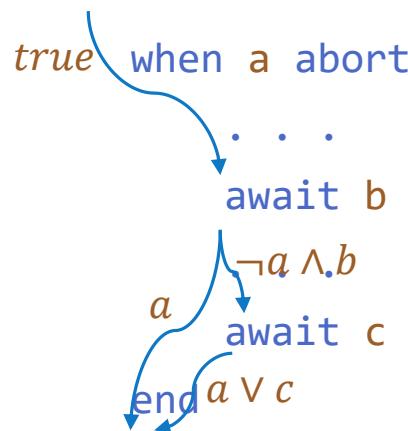
```
uint32_t main (int32_t *a, int32_t *b, int32_t *c, uint32_t *pc_1,
                uint32_t *pc_3, uint32_t *pc_2) {
    if ( *pc_1 == 2 ) {
        *a = 0; // init
        *b = 0;
        *c = 0;
        *pc_2 = 7; // enter branches and terminate step
        *pc_3 = 9;
        *pc_1 = 18;
    }
    if ( *pc_2 == 6 ) { // left branch
        *a = (*a + 2);
        *pc_2 = 12; // remember there is more to do
    }
    if ( *pc_3 == 8 ) { // right branch
        *b = (2 * (*a + 1));
        *pc_3 = 9; // terminate right step
    }
    if ( *pc_2 == 12 ) { // left branch
        *c = (*b - *a);
        *pc_2 = 7; // now terminate left step
    }
    _BLECH_SWITCH_TO_NEXTSTEP(*pc_2); // bit-shifting magic
    _BLECH_SWITCH_TO_NEXTSTEP(*pc_3);
    _BLECH_SWITCH_TO_NEXTSTEP(*pc_1);
    return *pc_1; // 0 means no more reaction steps to do
}
```

# Blech

## Stopping a behaviour

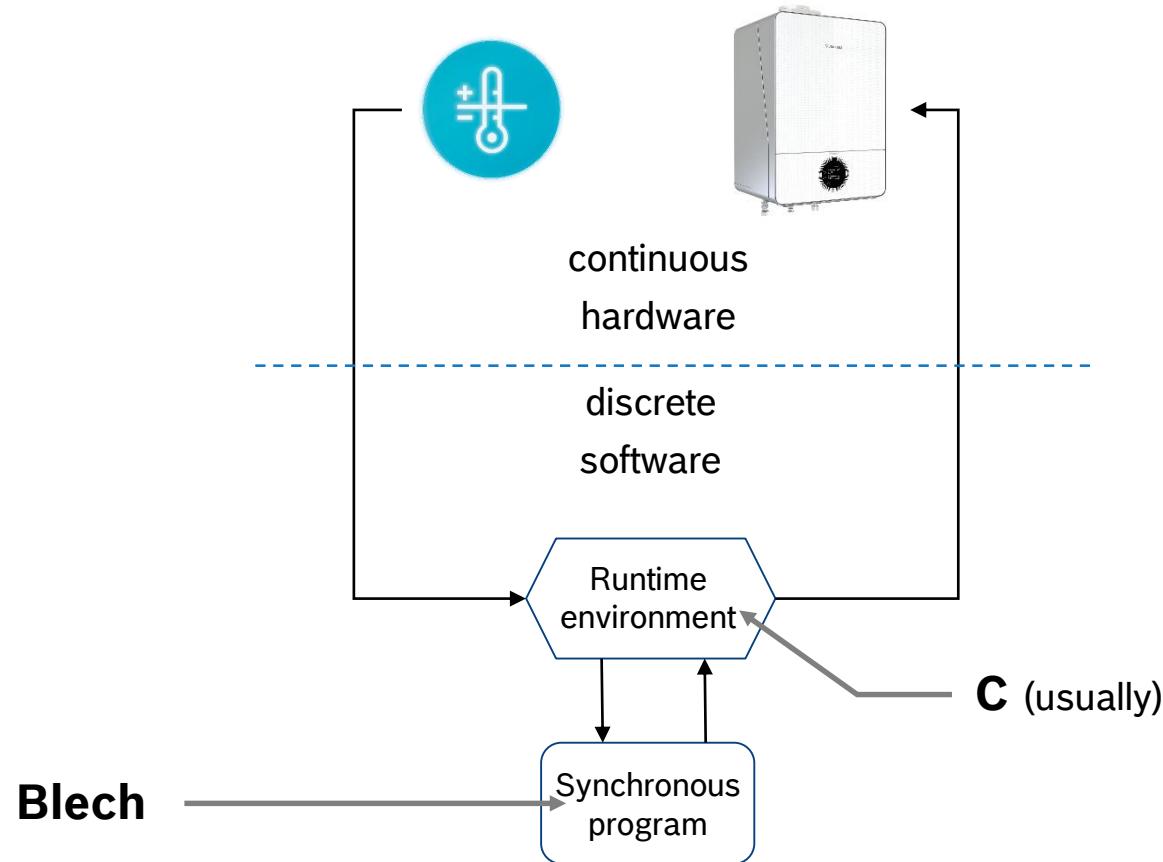
*/// Keep blinking until the user presses button 1*

```
activity Locked (pressedOne: bool) ()  
    when pressedOne abort  
        run Blink()  
    end  
end
```



# Abstract view of a reactive system

## Where do we use a synchronous language?



# C interoperability

## Calling Blech from a runtime

```
/* Main */

int main(int argc, const char * argv[])
{
    /* Create and initialize environment. */
    // ...
    /* Initialize blech. */
    blc_blech_acc_init();

    /* Sense, control, act loop */
    while (1) {

        /* Get and adapt sensor input from environment. */
        env_input_state_t env_input_state = env_read(env);
        // ...

        /* Run control reaction. */
        blc_blech_acc_tick(output_state.otherSpeed,
                            &output_state.egoSpeed,
                            &output_state.distance);

        /* Act on environment. */
        int hasCrashed = env_draw(env, &output_state);
        // ...

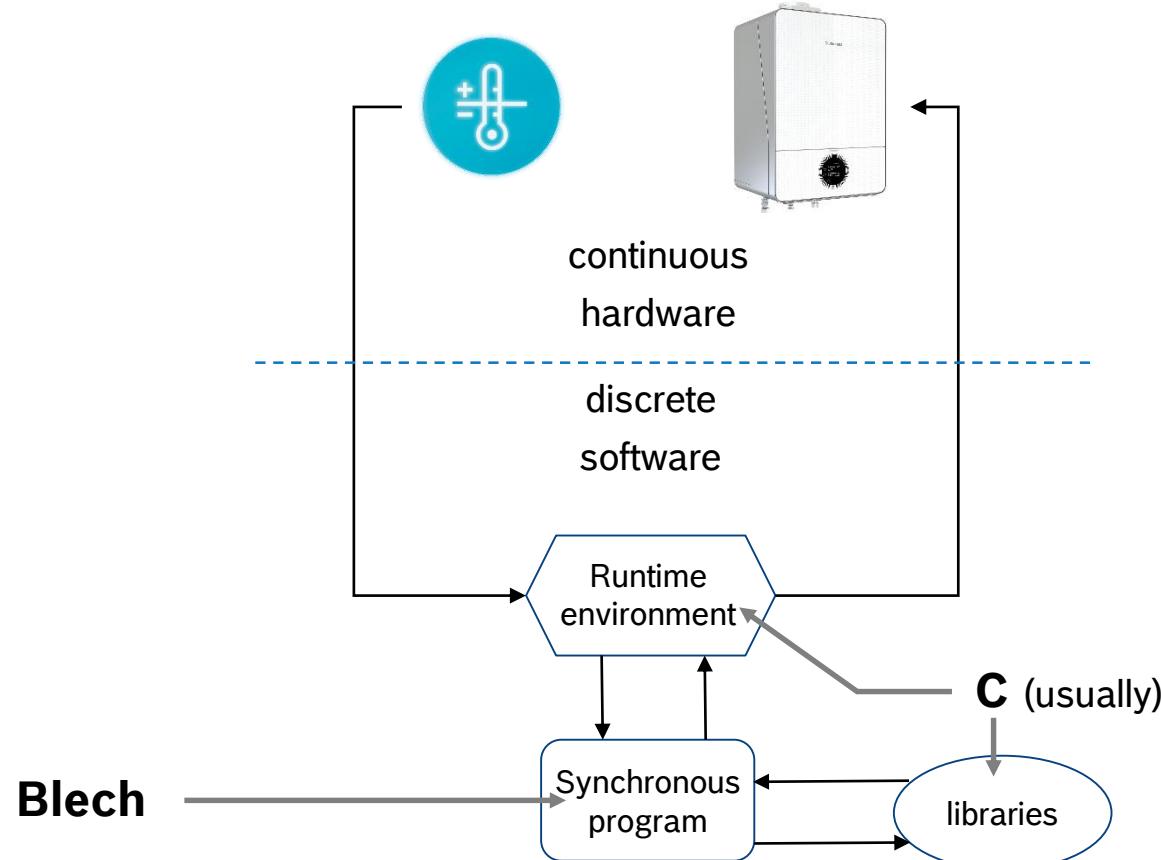
        /* Wait for next tick. */
        usleep(update_frequency);
    }

    /* Destroy environment. */
    env_destroy(env);

    return 0;
}
```

# Abstract view of a reactive system

## Where do we use a synchronous language?



# C interoperability

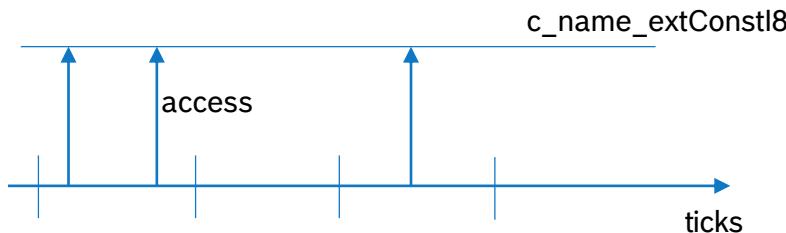
## External constants

C

```
#define c_name_extConstI8 8
```

### assumption

is constant throughout the whole runtime



### Blech

```
@[CConst(binding="c_name_extConstI8",
           header="my_exernals.h")]
extern const extConstI8: int32
```

### usage

```
function f ()
  let testI8 = extConstI8
  // ...
end
```

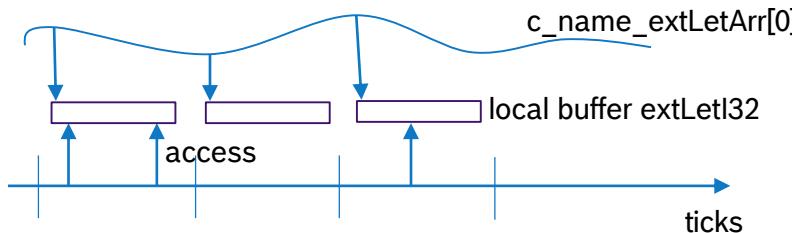
# C interoperability

## External volatile read-only memory

c

```
int c_name_extLetArr[8];
```

**assumption**  
is volatile



Blech

```
activity B ()  
  @[CInput(binding="c_name_extLetArr[0]",  
           header="my_externals.h")]  
  extern let extLetI32: int32  
  ...  
end
```

**usage** (multiple concurrent instances of B may run)

```
cobegin  
  run B ()  
with  
  run B ()  
end
```

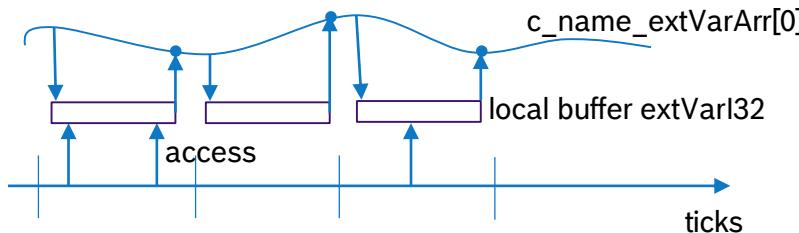
# C interoperability

## External volatile read-write memory

c

```
int c_name_extVarArr[8];
```

**assumption**  
is volatile



Blech

```
activity B ()  
  @[COutput(binding="c_name_extVarArr[0]",  
            header="my_externals.h")]  
  extern var extVarI32: int32  
  ...  
end
```

**usage** (B is a singleton now)

```
cobegin  
  run B ()  
with  
  run B ()  
end
```

error!

# C interoperability

## External (singleton) functions

C

```
uint8_t NRF24L01_spiIsReady (void)
{
    return (HAL_SPI_GetState(nrf24l01_init.hspi)
            == HAL_SPI_STATE_READY) ? 1 : 0;
}
```

Blech

```
@[CFunction(binding = "NRF24L01_spiIsReady",
             header = "nrf24l01.h")]
extern singleton function spiIsReady () returns bool
```

### assumption

singleton:

- function either reads a volatile value
- or has a side-effect on the environment

not singleton:

- re-entrant, side-effect free function

**usage** (spilsReady is declared to be a singleton)

```
cobegin
    await spiIsReady()
with
    await spiIsReady() ← error!
end
```

Find all details on the language as it is currently implemented at

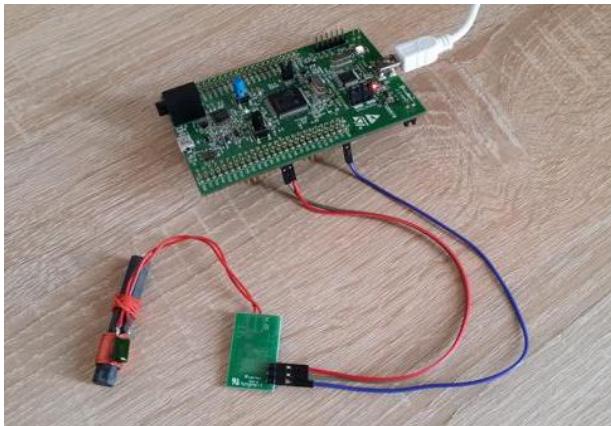
<https://www.blech-lang.org/docs/user-manual/>

If you find any mistakes or lack of clarity, please do notify us via Github issues.

# Overview

- ▶ Today's speaker
- ▶ Why is synchronous programming interesting for Bosch?
- ▶ Design goals
- ▶ Blech – as of now
- ▶ **Application examples**
- ▶ Outlook on planned features
- ▶ Additional remarks

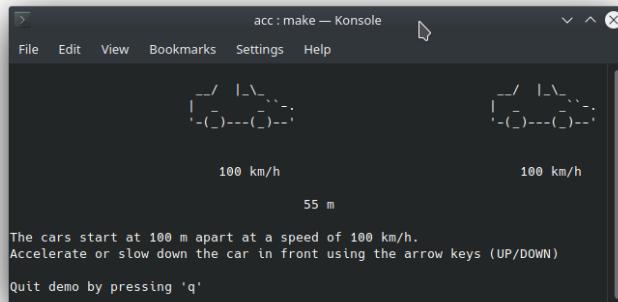
# Application examples



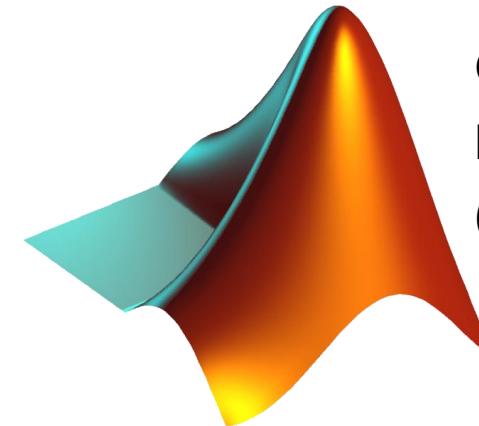
DCF 77 signal decoding  
bare metal



“Virtual lock”  
FreeRTOS + Mita



Homework: ACC  
Linux OS + ncurses



Controller development  
MATLAB/Simulink  
(S-function)

# Overview

- ▶ Today's speaker
- ▶ Why is synchronous programming interesting for Bosch?
- ▶ Design goals
- ▶ Blech – as of now
- ▶ Application examples
- ▶ **Outlook on planned features**
- ▶ Additional remarks

# Outlook on planned features

## What else should be possible with Blech?

### Mechanisms

- Parallel programming with multiple clocks
- Event communication using signals

### Software Engineering

- Module system
- Immutable references

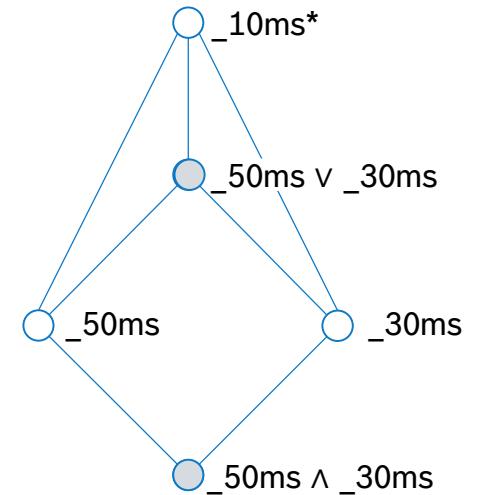
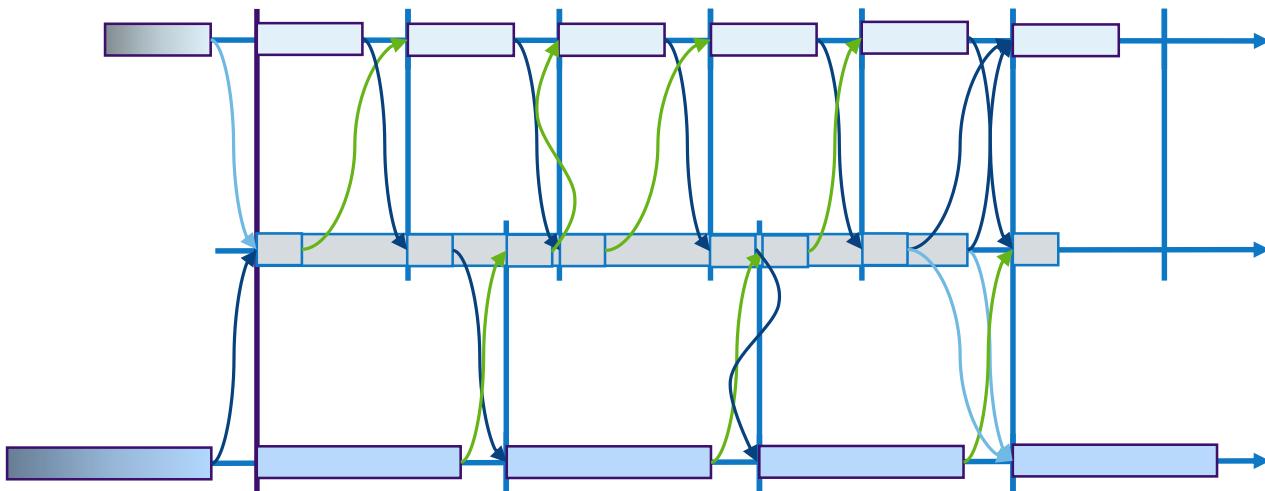
### Safety

- Physical dimensions
- Safe code generation

# Mechanisms

## Parallel programming with multiple clocks

- ▶ Communicating tasks must have related clocks
- ▶ Communication is done by sampling according to logical execution time
- ▶ Deterministic, consistent, compositional, real-time capable



# Mechanisms

## Communicating events with signals

```
activity Signalling()
  var finished: signal

  cobegin
    run anActivity()
    emit finished
  with
    repeat
      ...
      await true
    until finished end
    ...
  end
end
```

### Signal

- ▶ Presence flag
- ▶ Optional payload
- ▶ Only present in emitting time step
- ▶ Automatically absent after reaction

# Software engineering

## Module system

- ▶ Decompose code into separately compiled units: “modules” (do not confuse with Esterel modules!)
- ▶ Modules must export types, activities or functions that should be used by their clients (API, information hiding)
- ▶ Interfaces must take causality information into account
- ▶ Module system translates names to unique C identifiers (everything is globally visible in C)

# Safety

## Physical dimensions

```
unit m
unit s

var length: float32[m]
var duration: float32[s]

length = 2 * length // ok
length = 2 + length // error!

let speed = length / duration // ok
let nonsense = length + duration // error
```

- ▶ The physical dimension are part of the data type
- ▶ Machine data types prevent arithmetic operations on incompatible types
- ▶ Physical dimensions prevent arithmetic operations which do not make sense (cf. homework code)

# Safety

## Safe code generation

```
let a: [7]float32 = {...}  
...  
let x = a[i] // ok, provided i >= 0, i <= 6
```

### Debug code generation

```
float x;  
if(i >= 0 && i <= 6) {  
    x = a[i];  
} else {  
    haltWithDebugInfo();  
}
```

### Release code generation

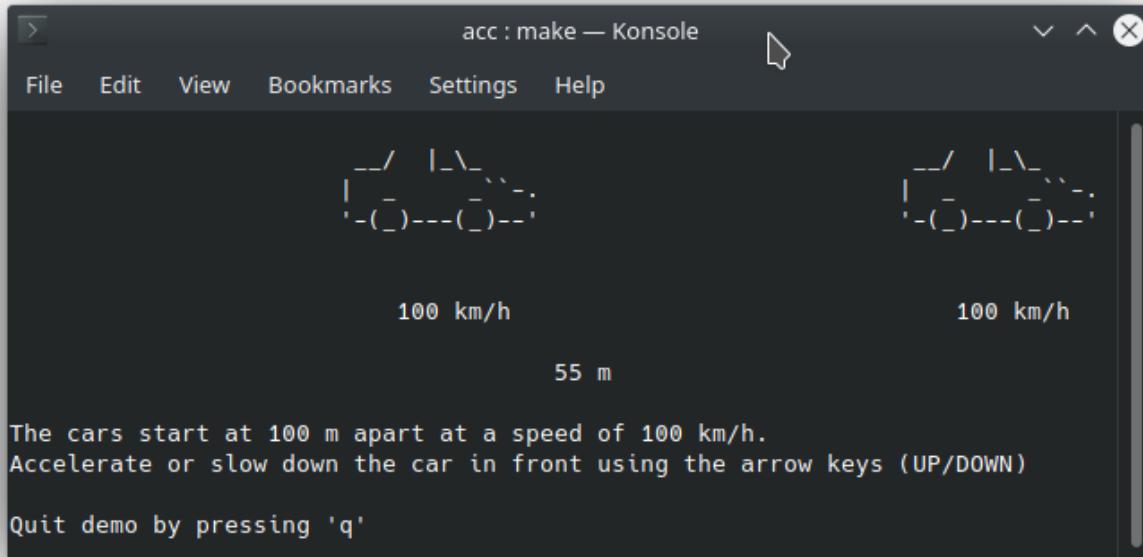
```
float x;  
if(i >= 0) {  
    if (i <= 6) {  
        x = a[i];  
    } else {  
        x = a[6];  
    }  
} else {  
    x = a[0];  
}
```

# Overview

- ▶ Today's speaker
- ▶ Why is synchronous programming interesting for Bosch?
- ▶ Design goals
- ▶ Blech – as of now
- ▶ Application examples
- ▶ Outlook on planned features
- ▶ **Additional remarks**

# Homework

## Adaptive Cruise Control



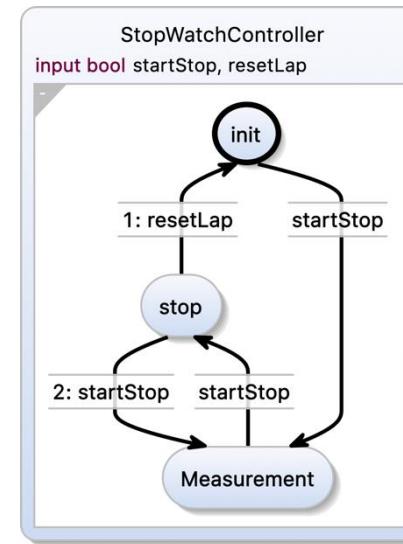
# Bachelor / Master Thesis

## Extraction of mode diagrams from Blech

Get in touch with Prof. von Hanxleden

```
activity StopWatchController (startStop: bool, resetLap: bool)
    (display: Display)

    var totalTime: int32
    var lastLap: int32
    repeat
        totalTime = 0      // State init
        lastLap = 0
        writeTicksToDisplay(totalTime)(display)
        await startStop   // Transition init -> run
        repeat
            cobegin weak
                await startStop
            with weak
                run Measurement(resetLap)
                    (totalTime, lastLap, display)
            end
            // State stop, show total time and wait
            writeTicksToDisplay(totalTime)(display)
            await startStop or resetLap
            // Run again if only startStop was pressed
        until resetLap end // Back to init if
                           // resetLap was pressed
    end
end
```



# Where do I get Blech? or how do I participate?

All info is available at [www.blech-lang.org](http://www.blech-lang.org)

Blech is **open source!** Driven by Bosch CR.

Try Blech right now, start with [tutorials and other examples](#). Why not write a [blog post](#) about your experience?

Participate in discussions and give feedback on language design  

Actively shape Blech by contributing to the [compiler](#), [tooling](#) or [documentation](#)

Let's collaborate on product software, an evaluation prototype, a student thesis or internship

We happily give a talk for your developers or managers or organise a hands-on tutorial

The Blech team is open for ideas and discussions

# Where do I get Blech? or how do I participate?

Get in touch with us:

[Friedrich.Gretz@de.bosch.com](mailto:Friedrich.Gretz@de.bosch.com)

[Franz-Josef.Grosch@de.bosch.com](mailto:Franz-Josef.Grosch@de.bosch.com)