

Observing Functional Logic Computations

or:

 Sy: The Curry Object Observation System

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C^{oo}Sy — Lazy + Logic = Hard Debugging

Curry: a lazy functional logic language (extension of Haskell)

Lazy (demand-driven) evaluation complicates debugging:

- execution trace does not match program text
- some terms are not evaluated
- print statements (for testing) might change program execution



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More problems by logic programming features:

- non-deterministic computations with multiple results
- instantiation of logic variables influences computation order
- information about bindings is relevant



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COOSy:

- a relatively simple approach to help debugging
- extension of HOOD (Haskell observation debugger [Gill'01])



COOSy is easy to use:

1. Import module **Observe**
2. Observe computed value of some expression e by

(observe observeType Label e)

(many **observeTypes** are predefined; see later)

3. Start graphical COOSy interface, execute program,
look at *observation protocol*



CSy — An Example

The following program contains a bug:

```
max x y | x < y = y  
        | x > y = x
```

```
maxList = foldl max 0
```

```
main = maxList [1,7,3,2,6,7,8]
```

Evaluate main \rightsquigarrow *no solution*



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Evaluate `main` \rightsquigarrow *no solution*

First debugging approach: observe the list

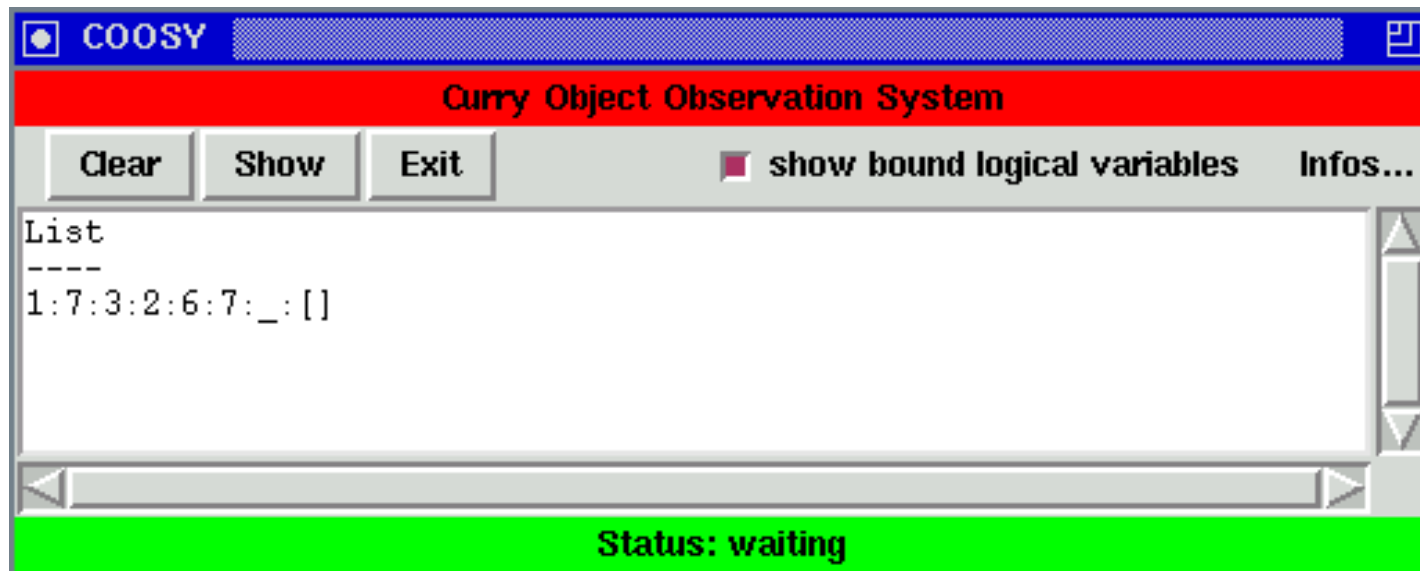
```
import Observe
```

```
...
```

```
main = maxList (observe (oList oInt) "List"
                    [1,7,3,2,6,7,8] )
```



Result of this example:



“_” \approx this element has not been evaluated



observe records uniquely numbered events:

Demand Event \approx a value is demanded:

Format:	Demand	argument	number	parent
Example:	Demand	1	24	23

Value Event \approx a value has been computed:

Format:	Value	value	arity	number	parent
Example:	Value	“4”	0	25	24

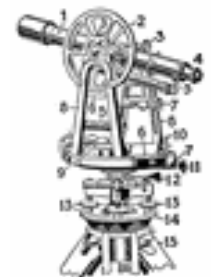


Fig. 1. 1. Tubus 2. 3. 4. 5. 6. 7. 8. 9. 10. 11. 12. 13. 14. 15.

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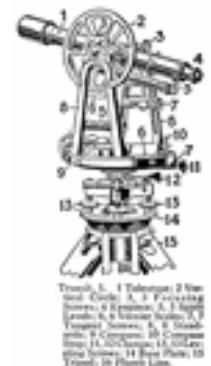
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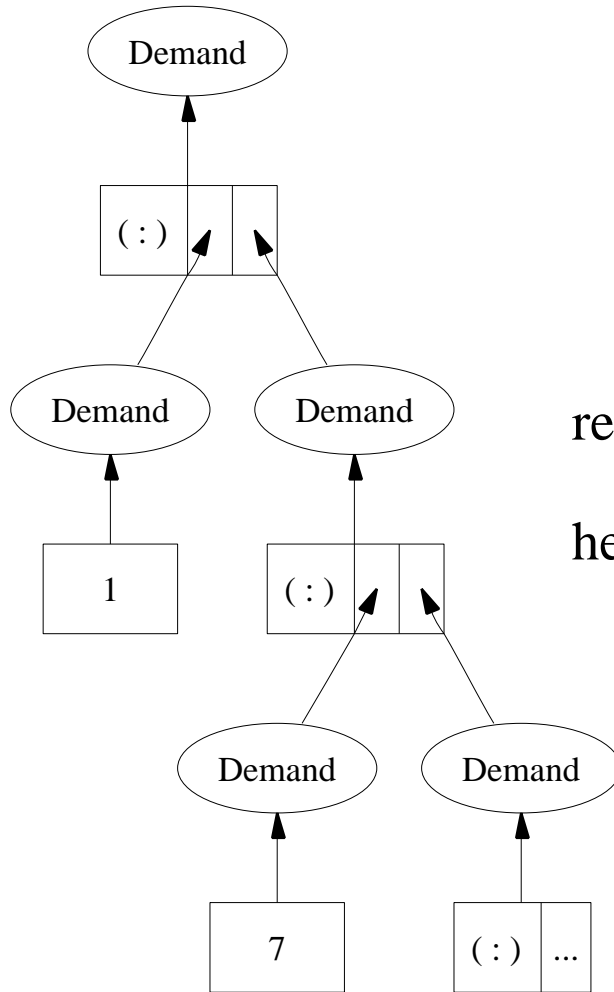
Format:	Value	value	arity	number	parent
Example:	Value	“4”	0	25	24

Chain of parent nodes \rightsquigarrow complete data structure:

- Value with reference r (arity > 0) but no Demand with parent r :
argument not demanded
- ! Demand with reference r but not Value with parent r :
value was demanded but not computed (failure, interrupt)

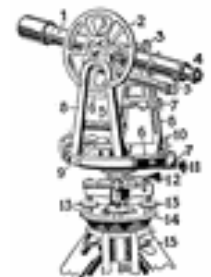


C^oSy — Reconstructing the Data Structure



reconstruct data structure from parent references

here: beginning of list $1 : 7 : [\dots]$



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COOSy — The Observables

One can observe

- data structures: standard observation types are derivable from general patterns

```
data Nat      = 0 | S Nat
oNat 0       = o0 "0" 0
oNat (S x)   = o1 oNat S "S" x
```



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- functions:

```
oMax      = oInt ~> oInt ~> oInt
maxList = foldl (observe oMax "max" max) 0
```

~> is predefined infix operator of module **Observe**



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- non-deterministic branches (see below)
- bindings of logic variables (see below)



Proceed with debugging the initial example:

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Observation protocol of a function:

argument/result pairs computed during program execution

```
{ ! _ -> !
, 7 7 -> !
, 7 6 -> 7
, 7 3 -> 7
, 1 7 -> 7
, 0 1 -> 1}
```



COSy — Non-Determinism

We want to observe also non-deterministic operations:

```
coin = 0
```

```
coin = S 0
```

```
main = plus 0 coin
```

```
plus 0      x = x
```

```
plus (S x) y = S (plus x y)
```



COOSy — Non-Determinism

We want to observe also non-deterministic operations:

```
coin = 0           plus 0       x = x
coin = S 0        plus (S x) y = S (plus x y)
main = plus 0 coin
main = (observe (oNat ~> oNat ~> oNat) " + " plus) 0 coin
```



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```

Desired observation:

```
+
---
{ 0 0 -> 0 }
{ 0 (S 0) -> S 0 }
```

But: current chaining of events provides not enough information



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Desired observation:

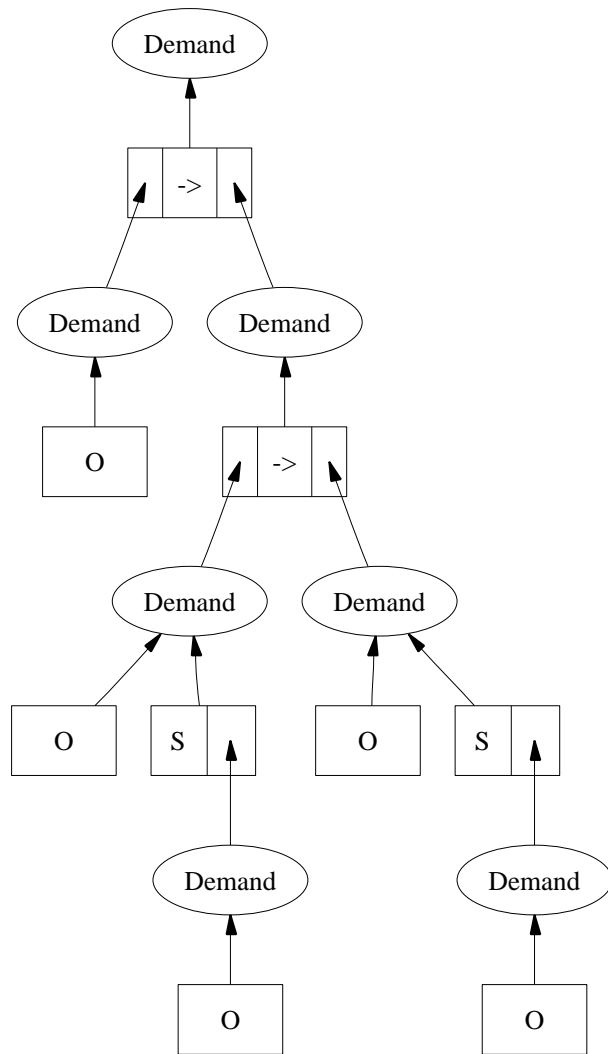
```
+
---
{ 0 0 -> 0 }
{ 0 (S 0) -> S 0 }
```

But: current chaining of events provides not enough information

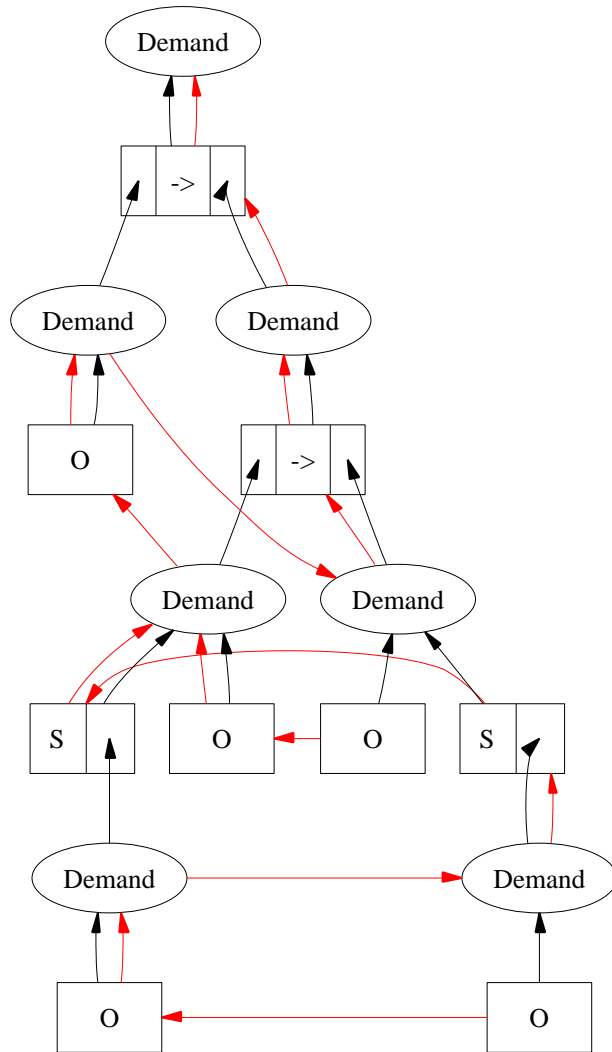
Predecessor: number of event occurred just before in same branch



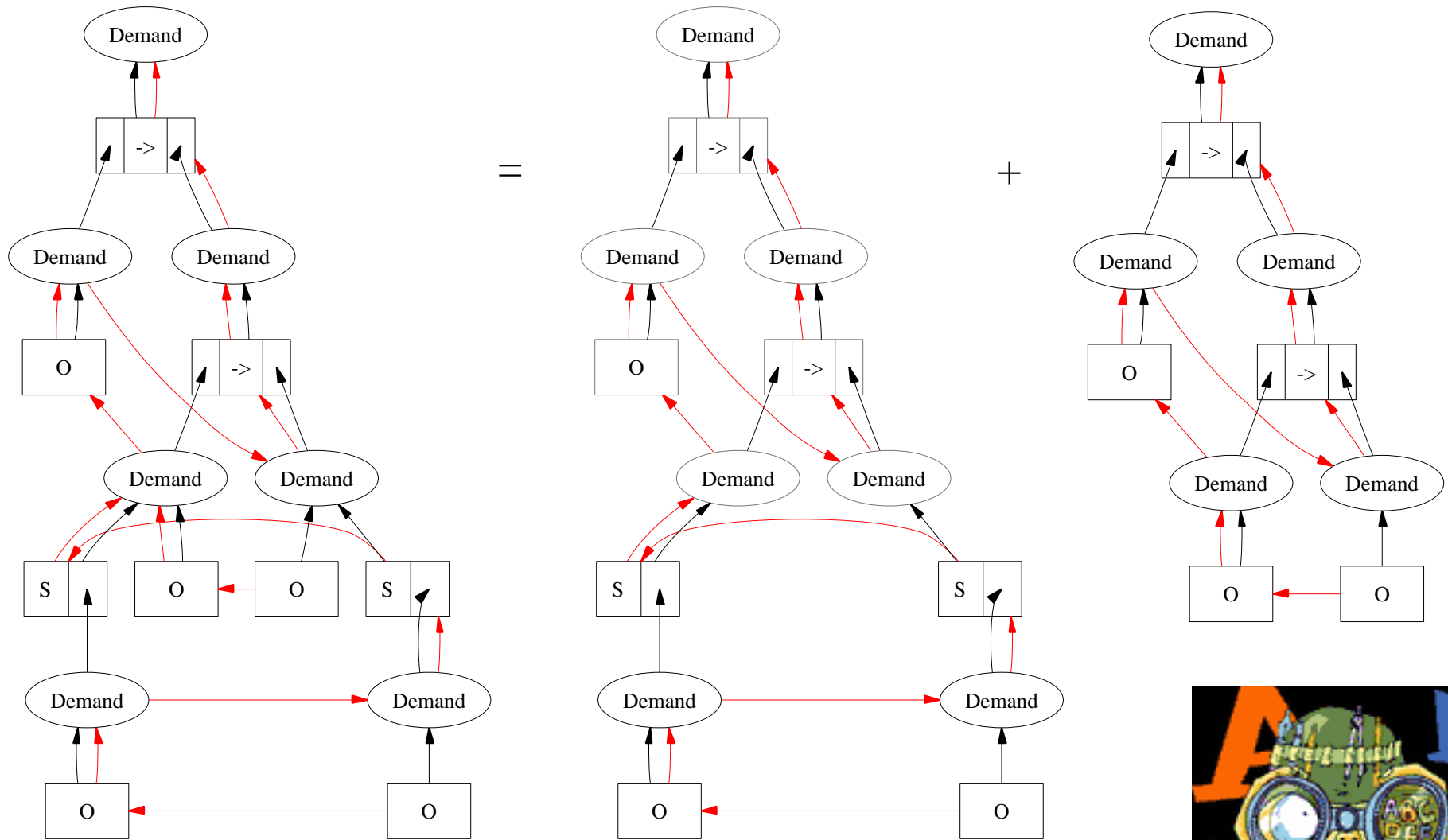
COSy — Predecessor Chaining



C^oSy — Predecessor Chaining



COSy — Predecessor Chaining



We want to observe also logic variables:

```
plus 0      x = x
```

```
plus (S x) y = S (plus x y)
```

```
main | plus          x y ::= S(S(S 0))  
     = (x,y)   where x,y free
```

~> four results: (0,3), (1,2), (2,1), (3,0)



We want to observe also logic variables:

```
plus 0      x = x
plus (S x) y = S (plus x y)
main | plus (observe oNat "x" x) y ::= S(S(S 0))
      = (x,y)   where x,y free
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↪ four results: (0,3), (1,2), (2,1), (3,0)

Observation for logic variable x:

```
(?/0)
(?/(S ?/0))
(?/(S ?/(S ?/0)))
(?/(S ?/(S ?/(S ?/0))))
```



C^oSy — How it Works

Problem: observation should not influence lazy evaluation

Solution: **observe** yields head constructor and further **observe** calls for arguments

Pattern: **observe** s (C $x_1 \dots x_n$) = C (**observe** 1 $x_1 \dots$ **observe** n x_n)



C^oSy — How it Works

Problem: observation should not influence lazy evaluation

Solution: `observe` yields head constructor and further `observe` calls for arguments

Pattern: `observe s (C x1 ... xn) = C (observe 1 x1 ... observe n xn)`

Problem: some terms are not evaluated

Solution: distinguish between Demand and Value events

Pattern: `observe x = record Demand >> compute hnf x >>= record Value`



C^oSy — How it Works

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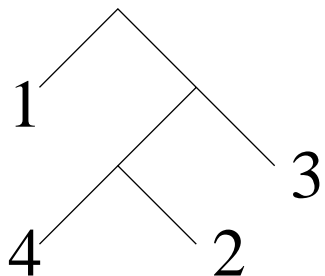
Solution: distinguish between Demand and Value events

Pattern: **observe** x = record Demand >> compute hnf x >>= record Value

Problem: associating subterms to computation branches

Solution: reference chain to predecessor of a branch

Pattern? such information is usually not accessible!



computation order can be arbitrary
(e.g., fair computation of all branches)



C^oSy — Computing Predecessor Chain

First solution: `observe` contains a logic variable as further parameter

- bindings visible in complete computation branch
- different computation branches — different bindings



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Remaining problem: when should this logic variable be bound?

- multiple bindings \rightsquigarrow failure of computation
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Solution:

- variable will be bound to an open-ended list, every step records reference at the end (by instantiation)
- find end of list by (unsafe) “test of logic variable” (`isVar`)



COSy — Observing Bindings

Final problem:

how can we observe concrete bindings of logic variables?

bindings are performed later (in other parts of the programs)



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how can we observe concrete bindings of logic variables?

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Solution:

create a concurrent constraint:

this constraint suspends on logic variable and is activated on binding

Pattern:

```
if isVar x then
  spawnConstraint (x ::= observe oNat "x" x)
else ...
```



COOSy Conclusions

COOSy extends Haskell's HOOD to include

- non-determinism
- logic variables



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All desired properties of HOOD are still valid:

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Advantages for debugging functional logic programs:

- no difficult requirements on underlying implementation:
functions used in module **Observe**: standard or easy to provide
- graphical interface \Rightarrow easy to use
- debug “batch” programs (e.g., web applications)



Automatic generation of **observe** types ✓

Distribute tool and test its usability

Printing of partial information in parallel to execution of observed program

