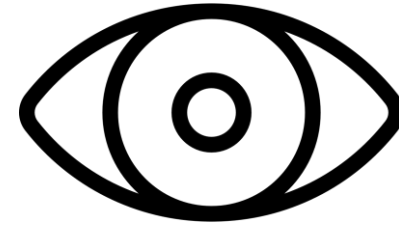


From Lustre to Graphical Models and SCCharts

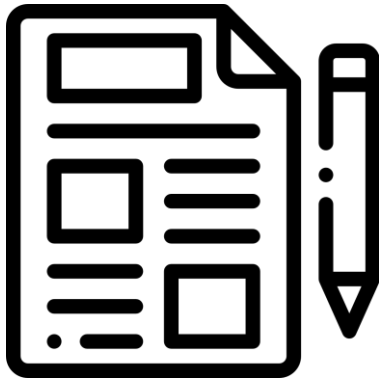
Lena Grimm
Kiel University



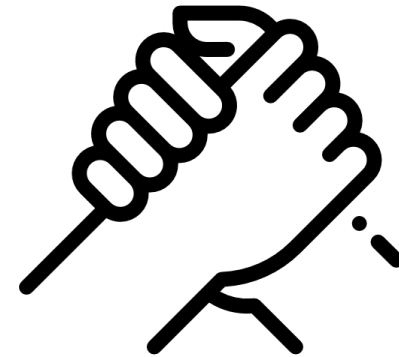
Motivation



Graphical Model



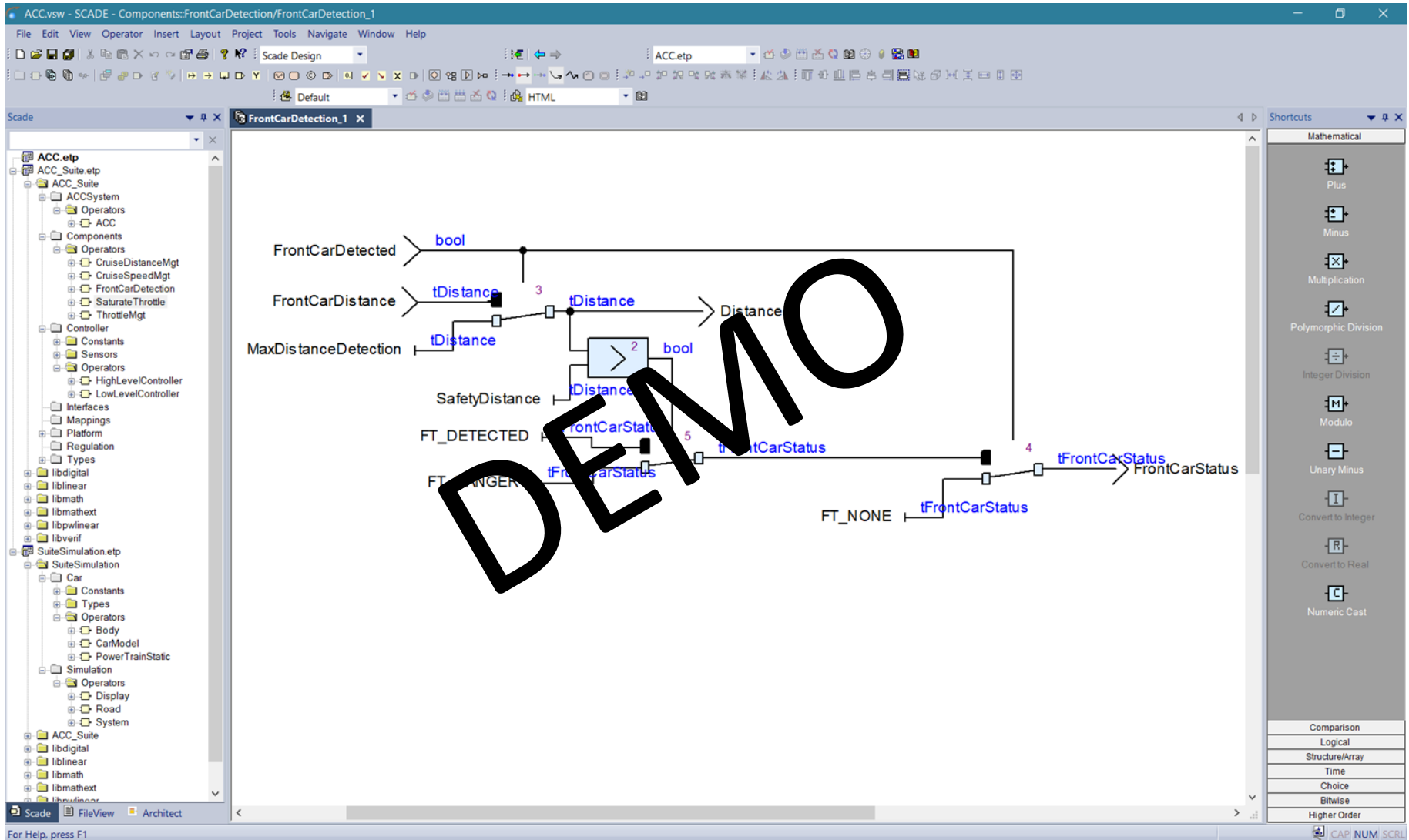
Validation
with SCCharts

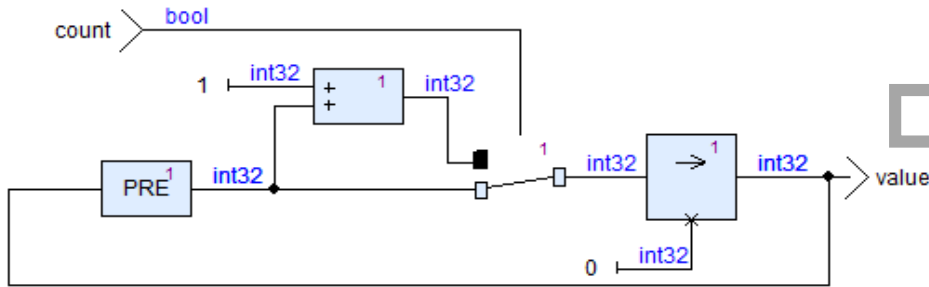


Sequential
Constructiveness



SCADE





```

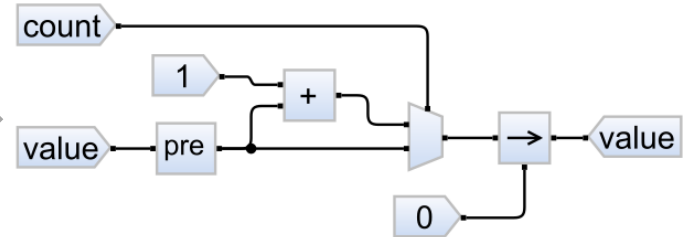
L1 = count;
value = L7;
L2 = 1;
L3 = pre L7;
L4 = L2 + L3;
L5 = if L1 then (L4) else (L3);
L6 = 0;
L7 = (L6) -> (L5);

```

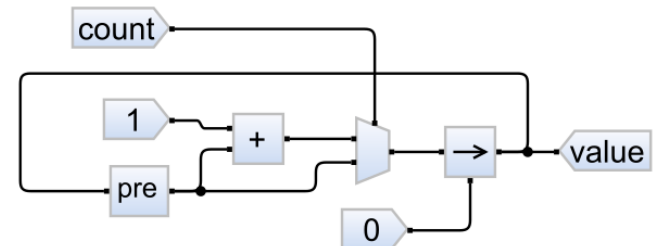
```

node Counter (count:bool)
  returns (value:int)
  var L1:bool;
  L2, L3, L4, L5, L6, L7:int;
  let
    L1 = count;
    value = L7;
    L2 = 1;
    L3 = pre L7;
    L4 = L2 + L3;
    L5 = if L1 then (L4) else (L3);
    L6 = 0;
    L7 = (L6) -> (L5);
  tel

```

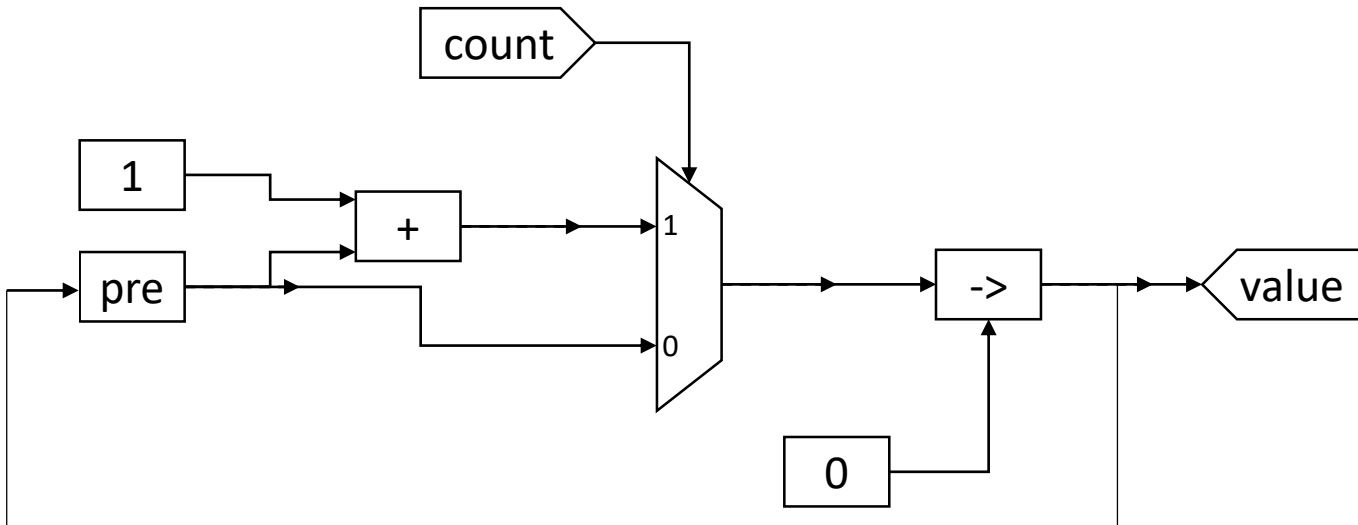
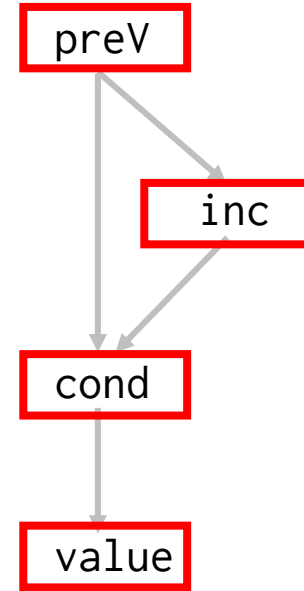
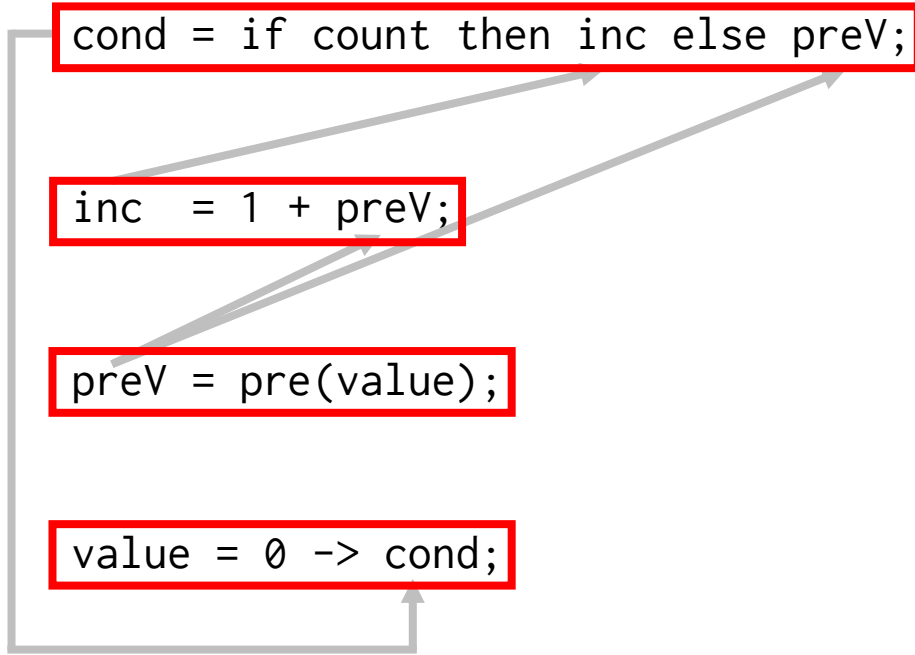


GENERATED



Input: count:bool

Output: value:int





KIELER Demo



```
node increment (count:bool) returns (value:int)
var cond:int;
    inc:int;
    preV:int;

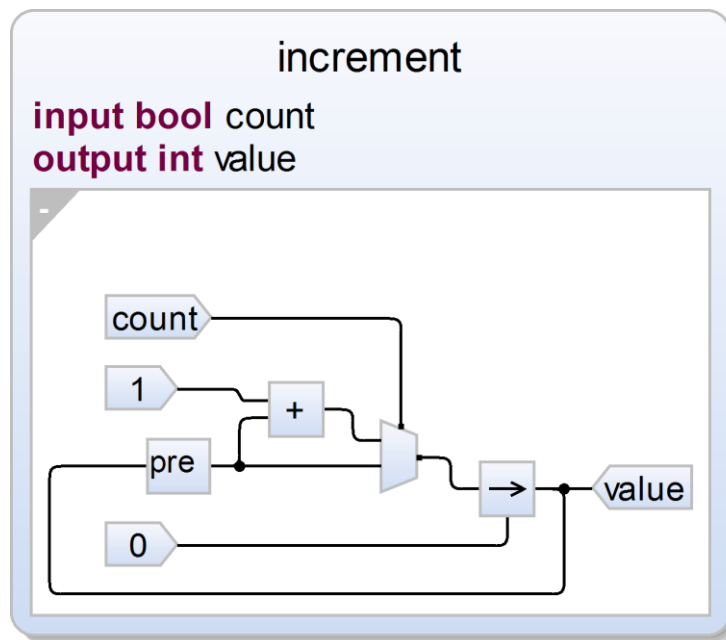
let
    cond = if count then inc else preV;
    inc = 1 + preV;
    preV = pre(value);
    value = 0 -> cond;
tel
```



```
scchart increment {
    input bool count
    output int value

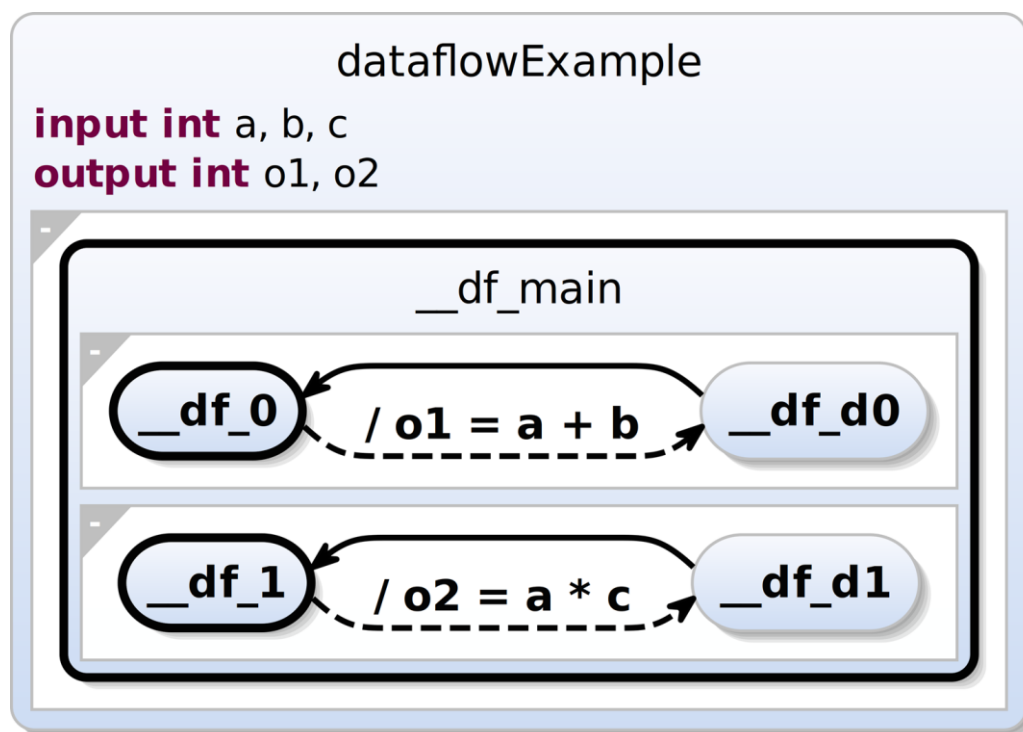
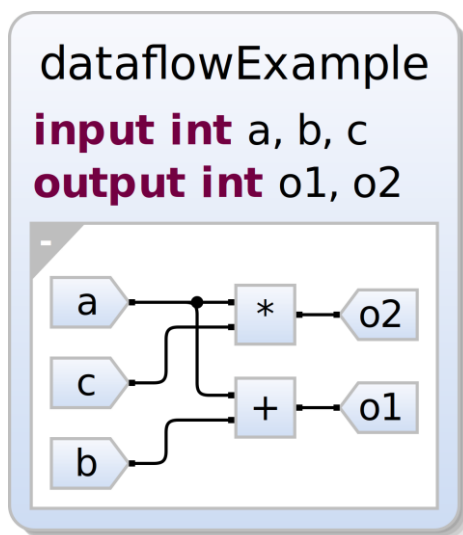
    dataflow {
        int cond, inc, preV

        cond = count ? inc : preV
        inc = 1 + preV
        preV = pre(value)
        value = 0 -> cond
    }
}
```



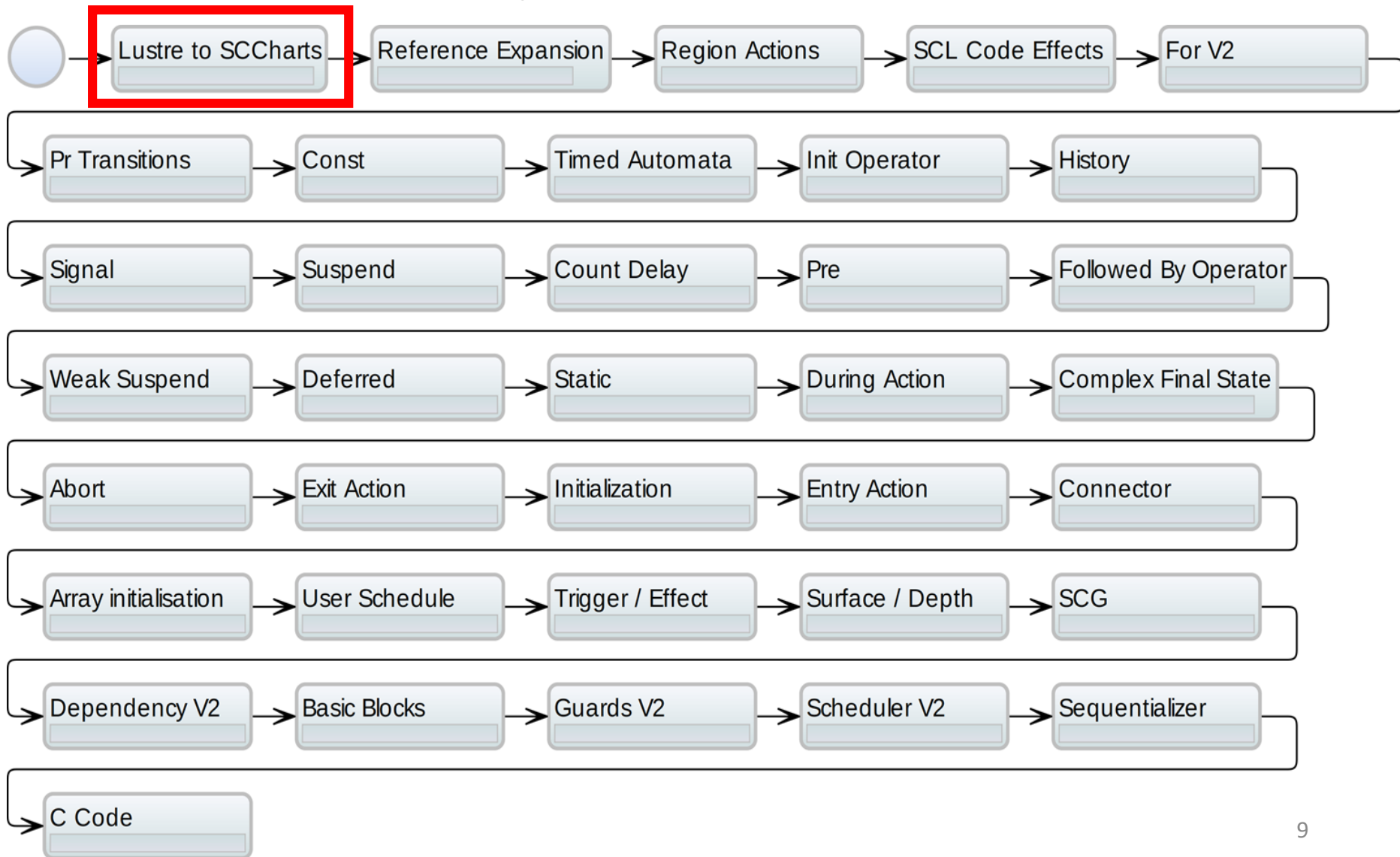


SCCharts Dataflow Semantics





KIELER Compilation Chain





Lustre Clock Calculus

Lustre

x	1	2	3	4	5	6	7	8	9
clk	true	false	true	false	false	true	false	true	true
x when clk	1		3			6		8	9



when, current?

SCCharts

x	1	2	3	4	5	6	7	8	9
clk	true	false	true	false	false	true	false	true	true
clk? x	1	1	3	3	3	6	6	8	9

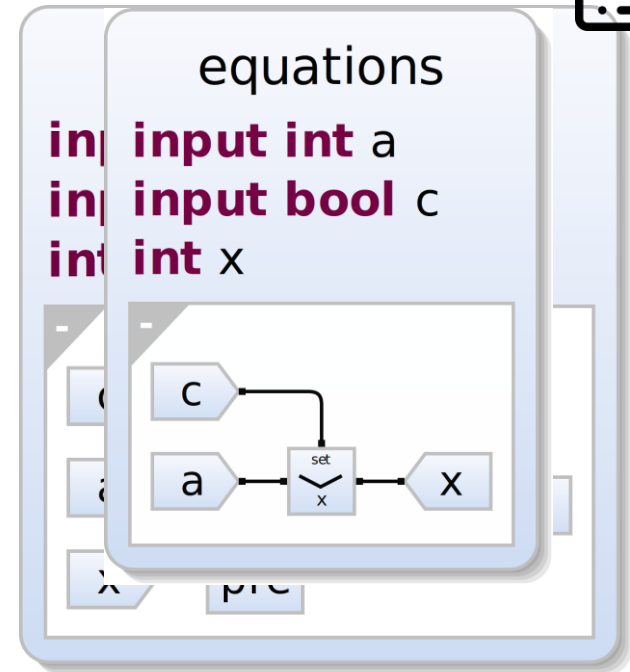
When Operator



```
node equations(a:int; c:bool)
  returns ();
```



```
var x:int when c;
let
  x = a when c;
tel.
```



```
scchart equations {
  input int a
  input bool c
  int x
  dataflow {
    x = c ? a : pre(x)
  }
}
```



Lustre When Operator

clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false
xClk = x when clk	true		false		true	false			true
y when xClk	true				false				false



When Operator with Variables I

clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false

xClk = clk? x	true	true	false	false	true	false	false	false	true
xClk? y	true	false	false	false	false	false	false	false	false



Hierarchical When

```
node equations(clk,x,y:bool)
  returns ();
```

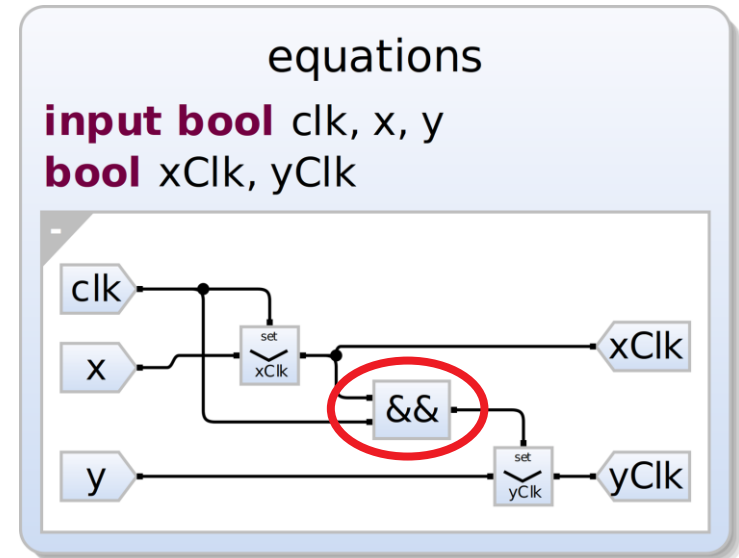
```
var xClk:bool when clk;
    yClk:bool when xClk;
```

```
let
```

```
    xClk = x when clk;
```

```
    yClk = y when xClk;
```

```
tel.
```



```
scchart equations {
  input bool clk, x, y
  bool xClk, yClk
```

```
  dataflow {
```

```
    xClk = clk ? x
```

```
    yClk = (xClk && clk) ? y
```

```
  }
```

```
}
```



When Operator with Variables II

clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false

$xClk = clk ? x$	true	true	false	false	true	false	false	false	true
$(clk \& \& xClk) ? y$	true	true	false	false	false	false	false	false	false





Current?

clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false

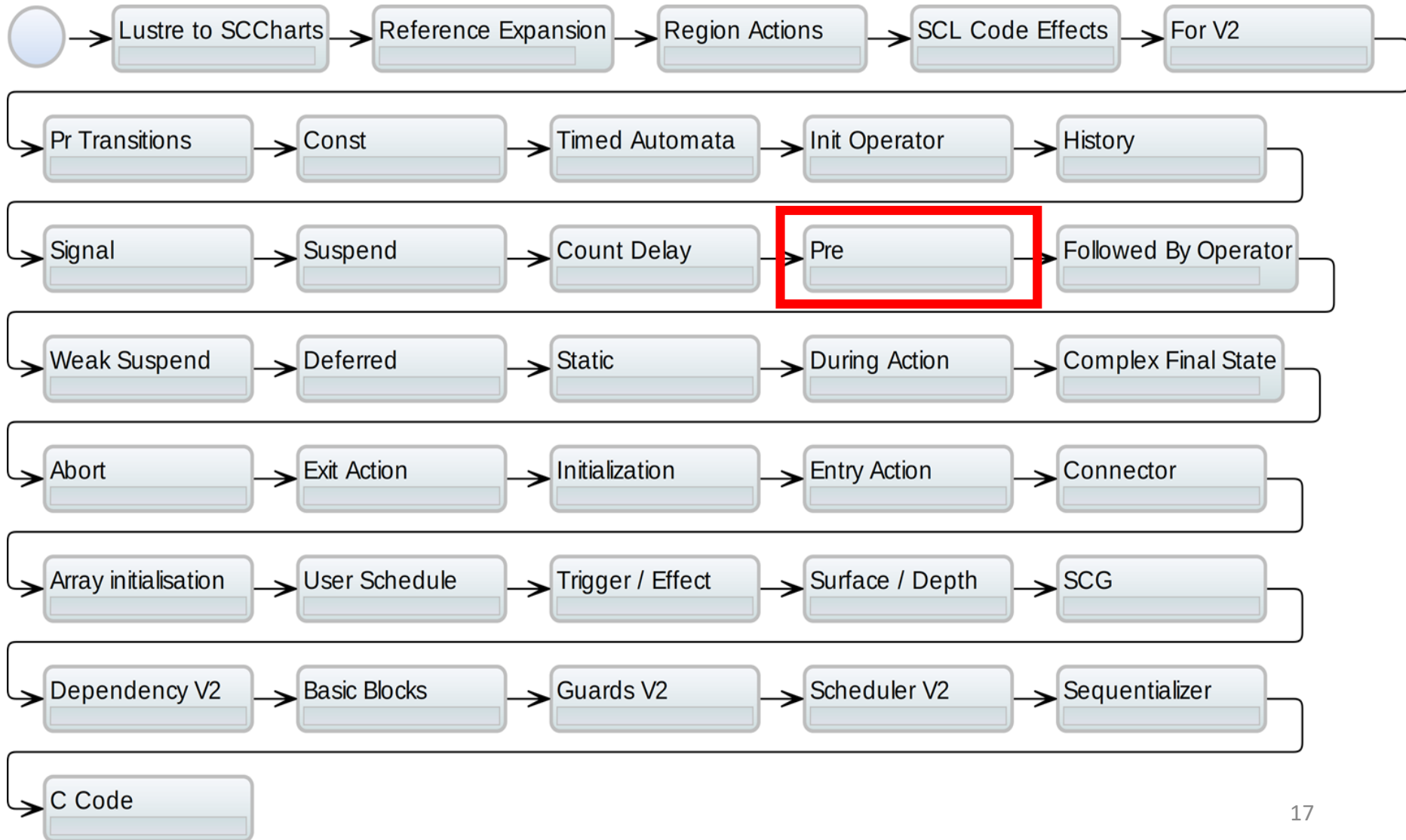
$xClk = clk ? x$	true	true	false	false	true	false	false	false	true
$(clk \& \& xClk) ? y$	true	true	false	false	false	false	false	false	false

Always implicit *current* through variables



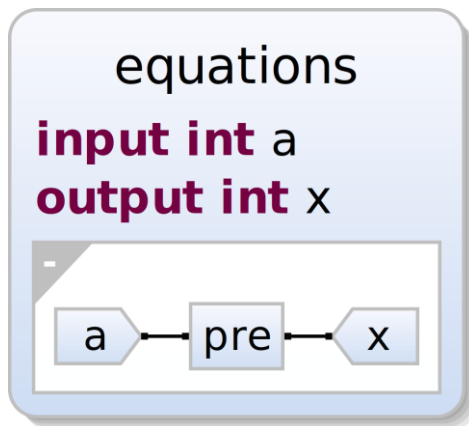


KIELER Compilation Chain

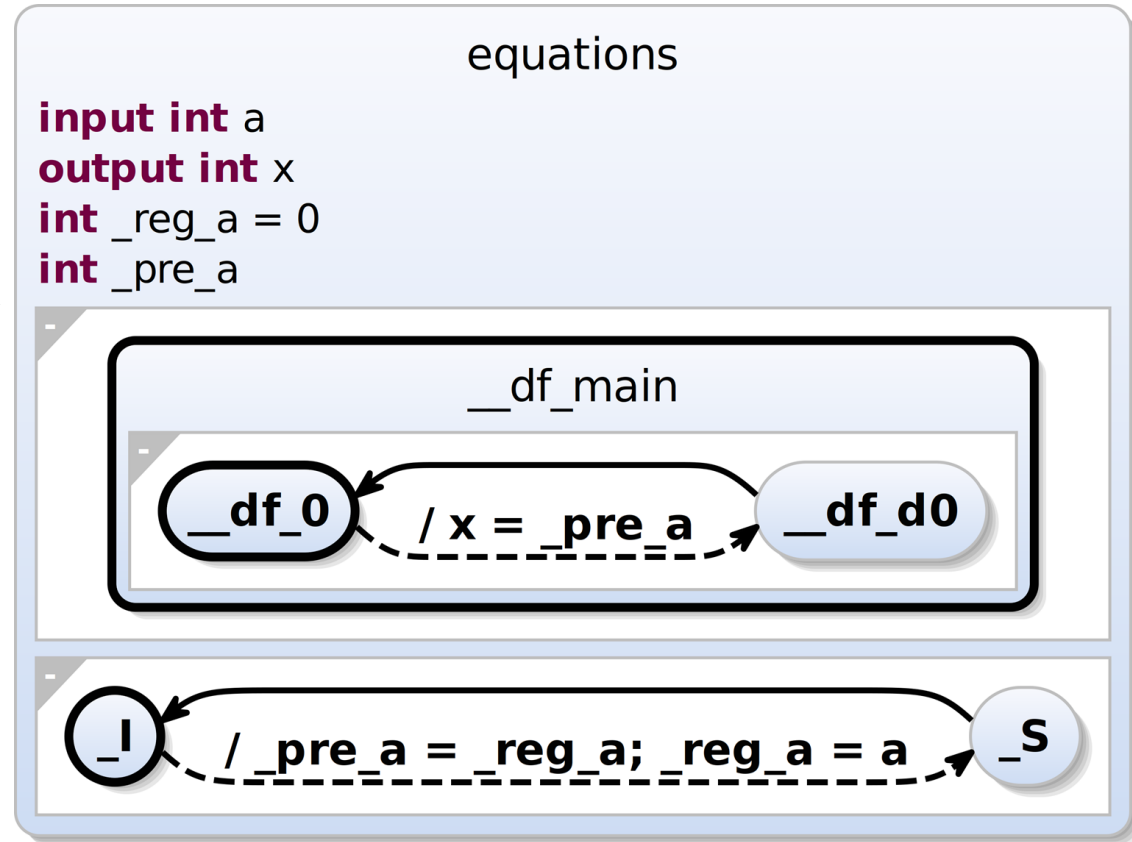


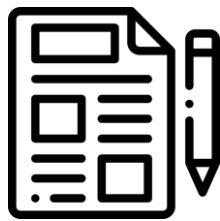


Pre Operator in SCCharts

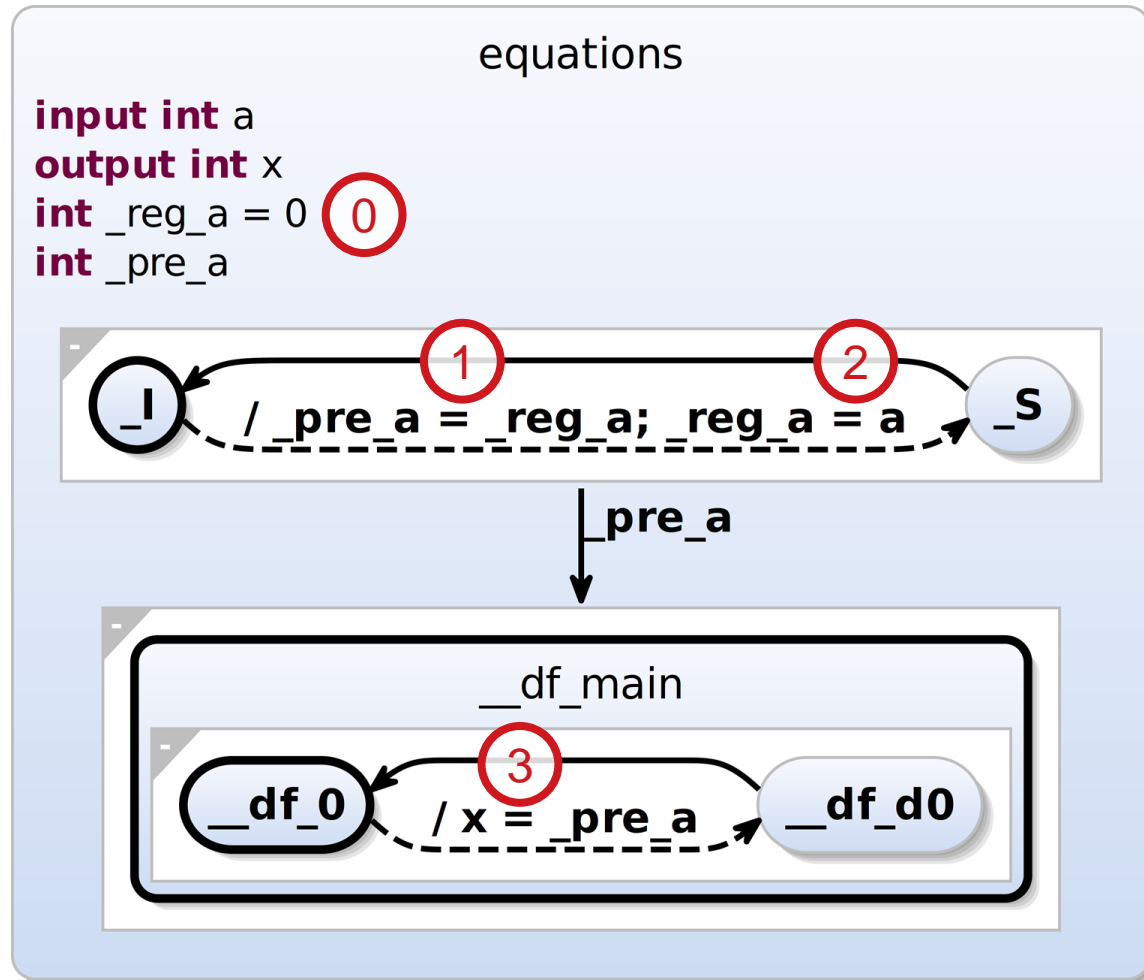


```
scchart equations {  
  input int a  
  output int x  
  
  dataflow {  
    x = pre(a)  
  }  
}
```





Pre Operator in SCCharts Induced Dataflow View





Lustre Pre Operator and Clocks

clk	true	false	false	true	true	false	true	false	true
x	1	2	3	4	5	6	7	8	9
xClk = x when clk	1			4	5		7		9
pxClk = pre(xClk)	nil			1	4		5		7
pre(pxClk)	nil			nil	1		4		5



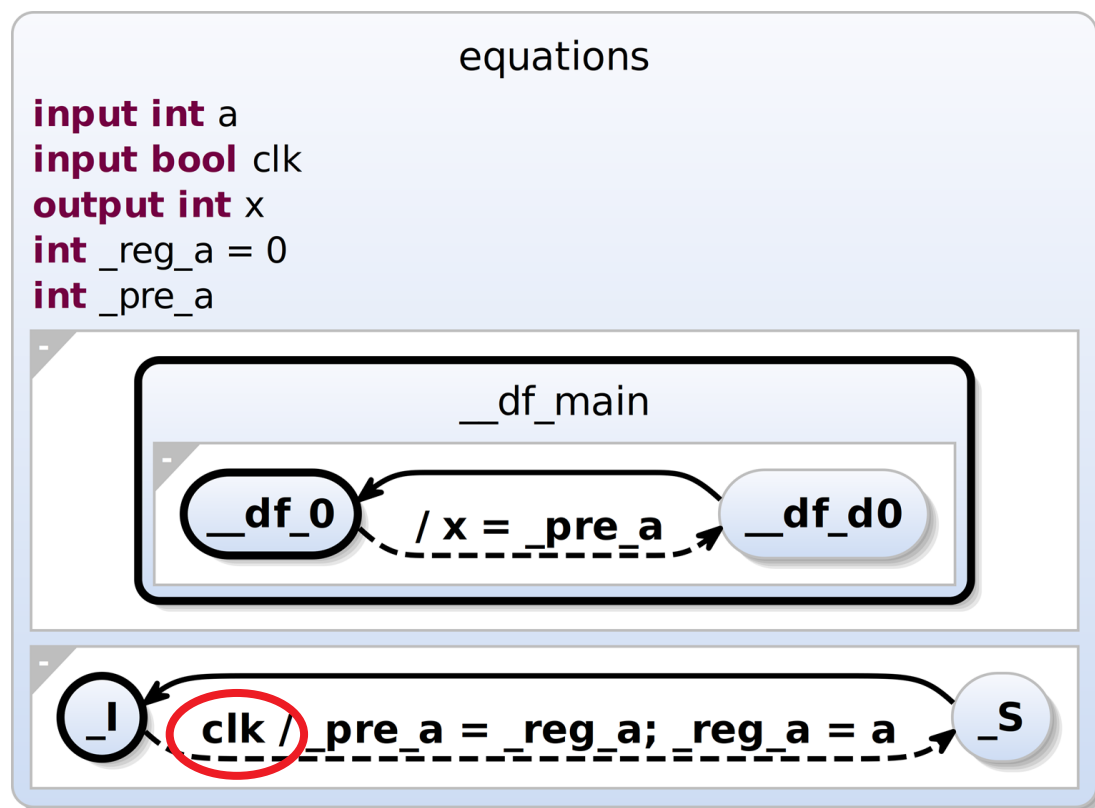
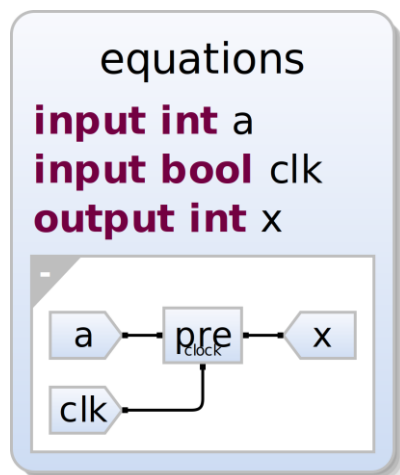
Pre Operator and Clocks with Variables

clk	true	false	false	true	true	false	true	false	true
x	1	2	3	4	5	6	7	8	9
xClk = clk? x	1	1	1	4	5	5	7	7	9
pxClk = pre(xClk)	nil	1	1	1	4	5	5	7	7
pre(pxClk)	nil	nil	1	1 nil	1	4	5 4	5	7 5

The table illustrates the behavior of the pre operator in a hardware description language. It shows the value of a variable 'x' over time, the value of 'xClk' (which is 'x' when 'clk' is true and the previous value of 'x' when 'clk' is false), the value of 'pxClk' (the previous value of 'xClk'), and the value of 'pre(pxClk)' (the previous value of 'pxClk'). Red arrows indicate the propagation of values from 'xClk' to 'pxClk' and from 'pxClk' to 'pre(pxClk)'. Red text highlights the values of 'pre(pxClk)' that are not nil, showing that the pre operator returns the value of 'pxClk' from the previous clock cycle.



Clocked Pre Operation in SCCharts



```
scchart equations {
  input int a
  input bool clk
  output int x
```

```
  dataflow {
    x = pre(a, clk)
  }
```

```
}
```



Clocked Pre Operator with Variables

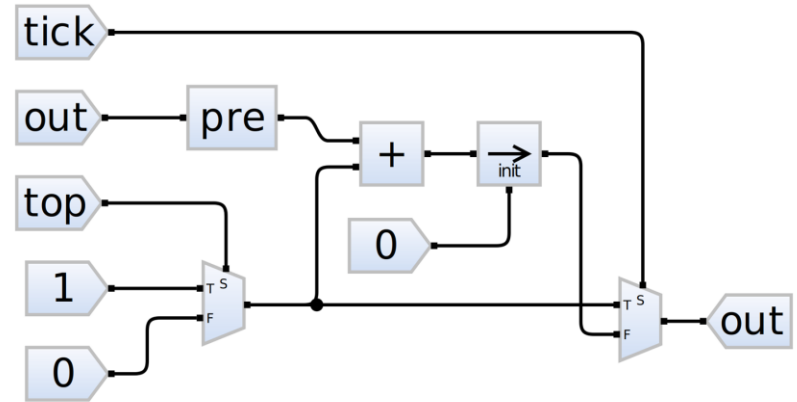
clk	true	false	false	true	true	false	true	false	true
x	1	2	3	4	5	6	7	8	9
xClk = clk? x	1	1	1	4	5	5	7	7	9
pxClk = pre(xClk)	nil	1	1	1	4	5	5	7	7
pre(pxClk)	nil	nil	1	nil	1	4	5	5	7



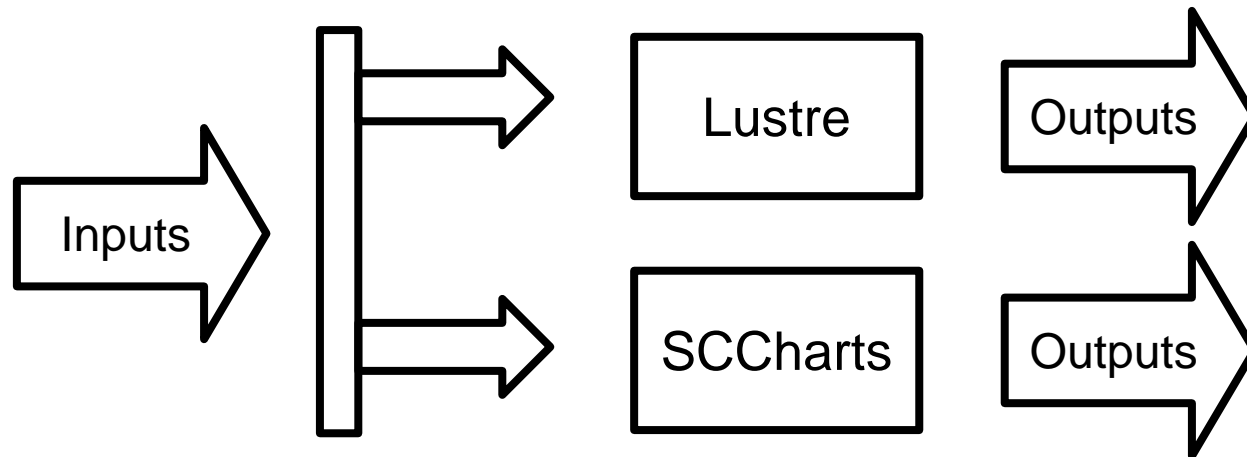
Model Recovery



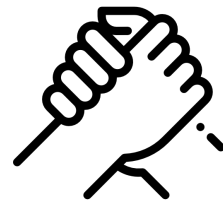
```
node counting(tick:bool;top:bool)
  returns (out:int);
var v:int;
let
  v = if top then 1 else 0;
  out = if tick
    then v
    else (0 -> pre out + v);
tel.
```



Behavior Preservation



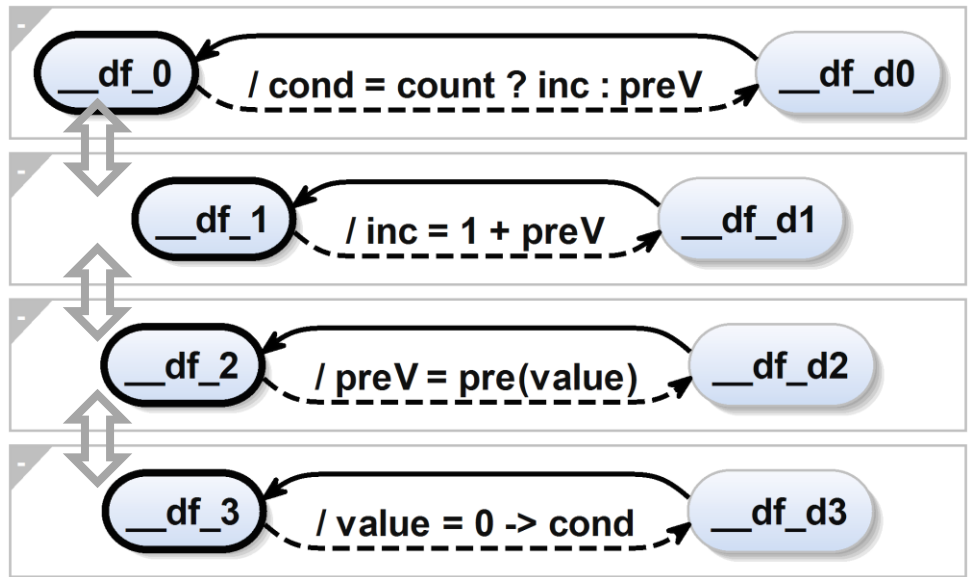
Sequential Constructiveness Concurrency



Initialize-Update-Read

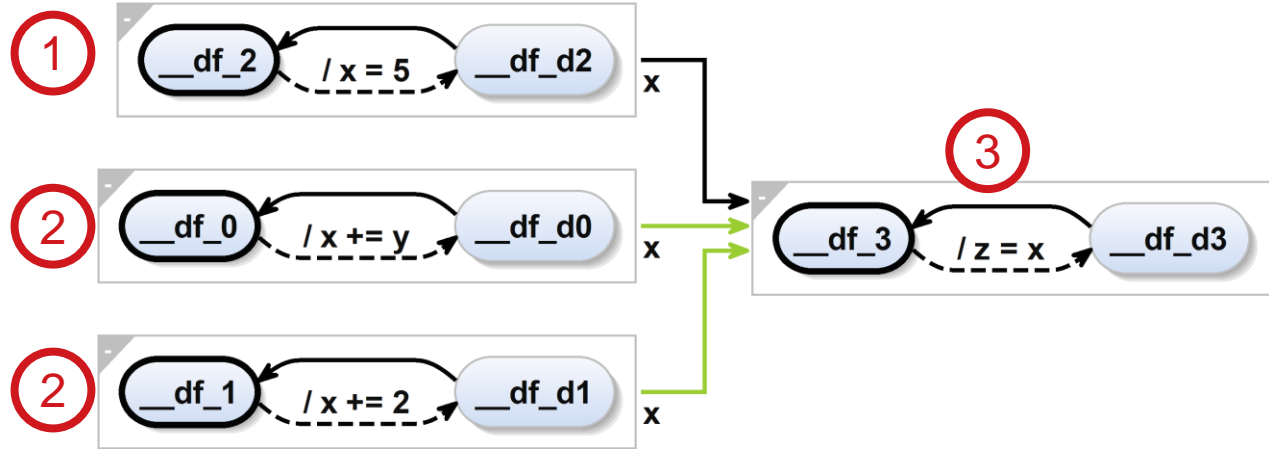
```
node increment (count:bool)
  returns (value:int)
var cond:int;
  inc:int;
  preV:int;

let
  cond = if count then inc else preV;
  inc = 1 + preV;
  preV = pre(value);
  value = 0 -> cond;
tel
```





...
x += y;
x += 2;
x = 5;
z = x;
...

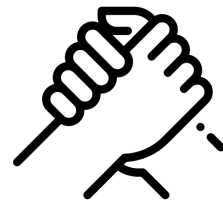


Conditioned Updates:

...
x = 5;
x += b? y;
x += c? 3;
...

Sequential Constructiveness

Sequentiality



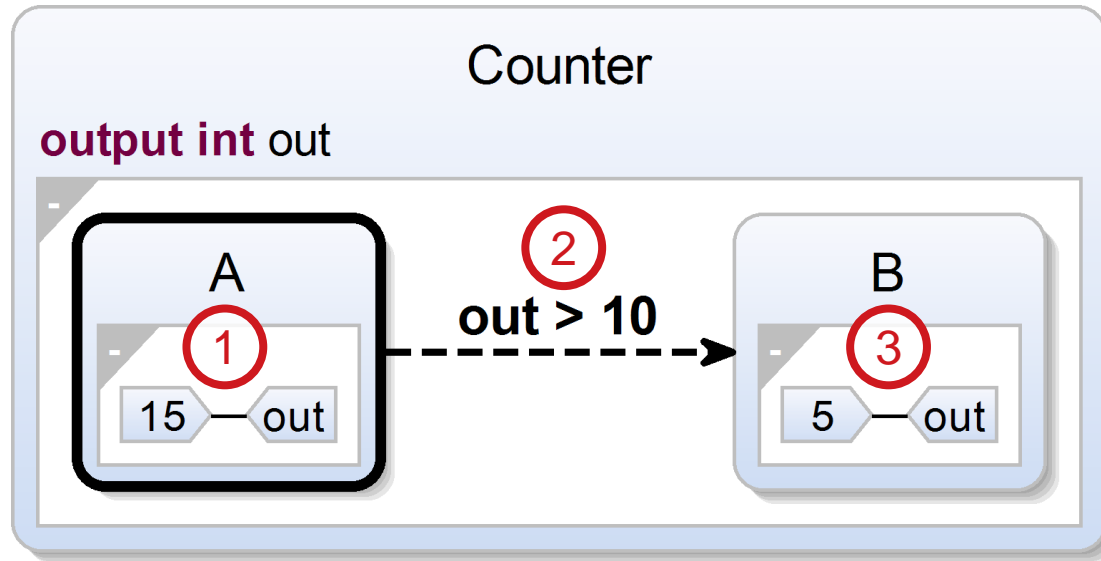
```
node simpleInc ()
  returns (value:int)
let
  value = 0 -> value + 1;
tel
```

No Register Variable needed



Sequential Constructiveness

Sequentiality with Automata



> Execute behavior of two states within one tick

Thank you!